

Catalan Clitics

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Les clitiques du catalan

RÉSUMÉ

Après une introduction à la grammaire catégorielle en ce qui concerne la structure basique de phrase, on offre un traitement du concord entre sujet et verbe parmi un formalisme qui comprend des types de logique de prédiqués. Cet appareil théorique est alors appliqué à l'analyse des clitiques pour laquelle on se sert de l'opération de "lifting". Des phénomènes tels que le doublage de clitique et le pro-drop de sujet sont traités dans le formalisme de la logique catégorielle. On rend compte du fait que la cliticisation se produit dans le domaine de la phrase en faisant référence aux domaines d'intensionnalité temporelle, lesquelles permettent d'exprimer ce que les clitiques peuvent franchir des verbes infinitivaux.

ABSTRACT

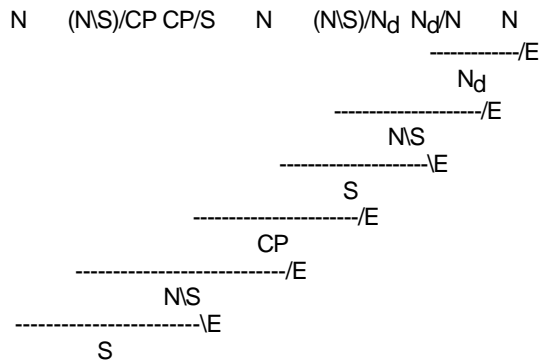
After a brief introduction to categorial grammar in relation to basic sentence structure, subject-verb agreement is implemented within the formalism with predicate-logical types. The apparatus is then applied to an analysis of clitics in terms of lifting. Phenomena such as clitic doubling and subject pro drop are treated in the general formalism of categorial logic. The clause-locality of cliticisation is captured by reference to temporally intensional domains in relation to which clitic-climbing over infinitival verbs is characterised.

CATEGORIAL GRAMMAR

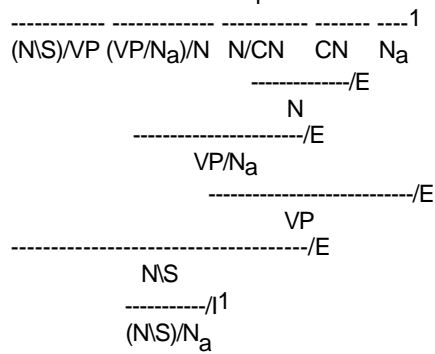
Let us regard a language as a system associating symbols and meanings; a model of a language will be built out of prosodic and semantic objects. A sign is a prosodic object/semantic object pair; a type is a set of signs. An (indexed) language model is a set of types (indexed by a set of type formulas), intended to match the association between symbols and meanings in various categories of the language observed. Prosodic and semantic objects are designated structurally by prosodic and semantic forms. An assignment is a prosodic form/semantic form/type formula triple. A formal language model is a set of assignments. A formal language model will be specified by a formal grammar, most simply a set of initial assignments (a lexicon) and a set of rules of formation under which the initial assignments are to be closed to generate the formal language model.

In categorial grammar, certain categories (names, statements, ...) are regarded as corresponding to signs which are primarily meaningful (or: complete). These categories are represented by basic types (N, S, ...). Further categories are identified as follows: if an expression

(3) El Joan pensa que el Pere parla de la Maria



(4) vol donar aquest llibre



PREDICATE-LOGICAL TYPES

Pure 'propositional' categorial grammar assumes unstructured atomic type formulas. This means that no relation between atomic types is expressed other than whether or not they are identical. Linguistic applications however demand a classificatory scheme which, for example, can indicate identity of major features despite difference of minor features in order that generalisations on the basis of major features can be captured. The straightforward generalisation of the formalism is to allow atomic type formulas to be Prolog-like first-order predicate-logical structures composed of feature constants, feature functions and type predicates, with feature variables being implicitly universally quantified at the outer level of a type formula. Then valid type inferences may be implemented by performing rules of inference with matching and unification on feature terms (see Morrill 1990a for this quantificational perspective on the role of unification).

Consider for instance the first conjugation as exemplified by the transitive verb *trobar*:

- (5) jo trobo nosaltres trobem
 tu trobes vosaltres trobeu
 ell, ella troba ells, elles troben

The value for agreement on the subject is constrained by the verb; there is no constraint on the object:

- (6) trobo - **find** := (N(1(sg))\S)/N(_)
 trobes - **find** := (N(2(sg))\S)/N(_)
 troba - **find** := (N(3(sg))\S)/N(_)
 trobem - **find** := (N(1(pl))\S)/N(_)
 trobeu - **find** := (N(2(pl))\S)/N(_)
 troben - **find** := (N(3(pl))\S)/N(_)

Proper names are categorized as follows.

- (7) el Joan - **john** := N(3(sg))
 la Maria - **mary** := N(3(sg))

Then derivation may proceed as illustrated in (8).

- (8) El Joan troba la Maria

 N(3(sg)) (N(3(sg))\S)/N(_)
 -----/E
 N(3(sg))\S
 -----\E
 S

The derivation induces a semantic term indicating how the semantics of the compound expression generated by the derivation is composed out of the semantics of the words from which it is formed. The semantic term assigned by the derivation above is (9a). On substitution of the lexical semantics this becomes (9b).

- (9) a. ((X_{troba} X_{la Maria}) X_{El Joan})
 b. ((**find mary**) **john**)

Similarly the lexical assignments in (10) induce the derivation in (11).

- (10) escriu - **write** := (N(3(sg))\S)/N_a(_)
 a - (λxx) := N_a(a)/N(a)
- (11) El Joan escriu a la Maria

 N(3(sg)) (N(3(sg))\S)/N_a(_) N_a(a)/N(a) N(3(sg))
 -----/E
 N_a(3(sg))
 -----/E
 N(3(sg))\S
 -----\E
 S

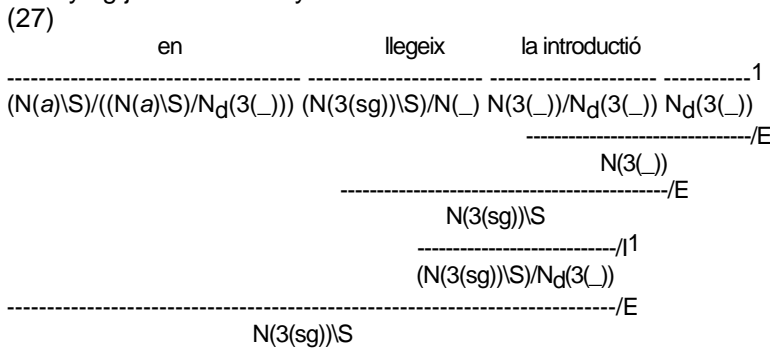
Since the semantics of the preposition is taken to be the identity function, substitution of the lexical semantics into the semantic term (12a) yields (12b) on simplification.

The clitic *en* is associated with a variety of functions (see e.g. Bartra 1987), amongst which is the satisfaction of a verb or noun *de* nominal complement requirement:

- (24) a. El Joan parla d'aquest llibre.
 'John talks about this book'
 b. El Joan en parla.
 'John talks about it'
- (25) a. El Joan llegeix la introducció d'aquest llibre.
 'John reads the introduction of this book'
 b. El Joan en llegeix la introducció.
 'John reads its introduction'

(26) *en* - **of-it** := (N(a)\S)/((N(a)\S)/N_d(3(_)))

Note that in (25b) the clitic attaches to a verb which is not the head of the complement position that the clitic binds: clitics are not limited to satisfying just the valency of their verb.



The clitic *hi* may perform locative adverbial modification, and *ho*, sentential complementation.

- (28) a. El Joan vagi a la biblioteca.
 John goes to the library
 b. El Joan hi vagi.
 'John goes there'
- (29) a. El Joan pensa que la Maria canta.
 'John thinks that Mary sings'
 b. El Joan ho pensa.
 'John thinks so'

(30) *hi* - **there** := (N(a)\S)/((N(a)\S))
ho - **so** := (N(a)\S)/((N(a)\S)/CP)

The second person clitics *et* (sg) and *us* (pl) can satisfy either a direct object or indirect object valency (dative clitics can also fulfill an 'ethic' role expressing someone to whom an act is beneficent; this aspect is not included here).

- (31) a. El Joan et renta.
 'John washes you(sg)'
 b. El Joan t' escriu.
 'John writes to you(sg)'

- (32) a. El Joan us rento.
 'John washes you(pl)'
 b. El Joan us escriu.
 'John writes to you(pl)'

The first person *em*(sg) and *ens*(pl) behave likewise.

- (33) a. El Joan em renta.
 'John washes me'
 b. El Joan m' escriu.
 'John writes to me'

- (34) a. El Joan ens renta
 'John washes us'
 b. El Joan ens escriu.
 'John writes to us'

In order to capture the polymorphism exhibited here a conjunctive type-constructor & is used (see Morrill 1990a). Note that the same semantics is given to cover the cases where the clitic binds direct and indirect object positions.

(35) $ens - (\lambda x(x\ 1(pl))) := (N(a)\S)/((N(a)\S)/(N(2(sg))\&N_a(2(sg))))$

The rules for the type-constructor are as follows:

(36) $\Gamma \Rightarrow \alpha : B \quad \Gamma \Rightarrow \alpha_1 : C$
 -----[&I], $\alpha = \alpha_1$
 $\Gamma \Rightarrow \alpha : B\&C$

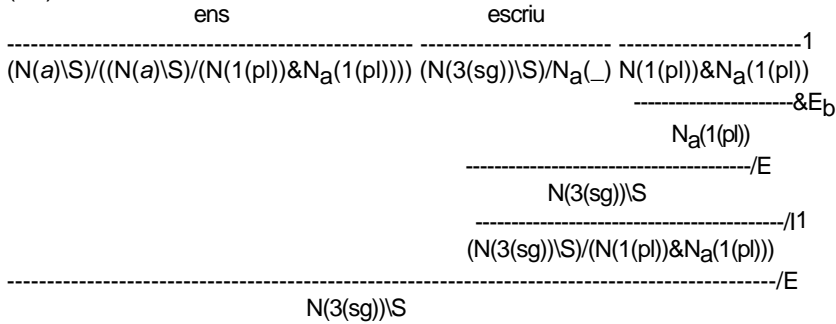
$\Gamma \Rightarrow \gamma : A\&B$ $\Gamma \Rightarrow \gamma : A\&B$
 -----[&E_a] -----[&E_b]
 $\Gamma \Rightarrow \gamma : A$ $\Gamma \Rightarrow \gamma : B$

The derivations in (37) and (38) illustrate how both a transitive verb and a prepositional verb inhabit the argument type of *ens*.

(37)

ens	renta
-----1	
(N(a)\S)/((N(a)\S)/(N(1(pl))\&N _a (1(pl)))) (N(3(sg))\S)/N() N(1(pl))\&N _a (1(pl))	
	-----&E _a
	N(1(pl))
	-----/E
	N(3(sg))\S
	-----/1
	(N(3(sg))\S)/(N(1(pl))\&N _a (1(pl)))
	-----/E
-----/E	
N(3(sg))\S	

(38)



The clitics *li* and *els* bind third person dative objects with which they optionally co-occur; in the event of such clitic doubling there must be agreement of number.

- (39) a. El Joan *li*/**els* escriu a la noia
 b. El Joan **li*/*els* escriu a les noies

In order to capture such possibilities intersection may again be used, but it is necessary for the semantics to be different depending on whether it is to be taken from a realised complement, or just the agreement features of the clitic.

- (40) *li* - $(\lambda x(x\ 3(sg)), x)$:=
 $((N(a)\S) \wedge ((N(a)\S)/N_a(3(sg)))) / ((N(a)\S)/N_a(3(sg)))$
els - $(\lambda x(x\ 3(pl)), x)$:=
 $((N(a)\S) \wedge ((N(a)\S)/N_a(3(pl)))) / ((N(a)\S)/N_a(3(pl)))$

The rules of use and proof for this semantically potent conjunction are semantically interpreted in terms of pairing and projection; the earlier, semantically impotent, conjunction was semantically interpreted by just identity.

- (41) $\Gamma \Rightarrow \alpha : B \quad \Gamma \Rightarrow \alpha_1 : C$
 -----[\wedge],
 $\Gamma \Rightarrow (\alpha, \alpha_1) : B \wedge C$

- | | |
|--|--|
| $\Gamma \Rightarrow \gamma : A \wedge B$ | $\Gamma \Rightarrow \gamma : A \wedge B$ |
| -----[$\wedge E_a$] | -----[$\wedge E_b$] |
| $\Gamma \Rightarrow (\pi_1 \gamma) : A$ | $\Gamma \Rightarrow (\pi_2 \gamma) : B$ |

The reduction laws for pairing and projection are given in (42).

- (42) $(\pi_1(\alpha, \beta)) = \alpha$
 $(\pi_2(\alpha, \beta)) = \beta$

Clitic doubling and non-clitic doubling derivations are as follows.

$$\begin{array}{c}
(43) \qquad \qquad \qquad \text{els} \qquad \qquad \qquad \text{escriu} \\
\hline
((N(a)S) \wedge ((N(a)S)/N_a(3pl))) / ((N(a)S)/N_a(3pl)) \ (N(3sg))S / N_a(L) \\
\hline
(N(3sg))S \wedge ((N(3sg))S)/N_a(3pl)) \\
\hline
\qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \wedge E_b \\
\qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad (N(3sg))S / N_a(3pl))
\end{array}$$

$$(44) \quad (\pi_2(x_{els} \ x_{escriu})) \\
\text{write}$$

$$\begin{array}{c}
(45) \qquad \qquad \qquad \text{els} \qquad \qquad \qquad \text{escriu} \\
\hline
((N(a)S) \wedge ((N(a)S)/N_a(3pl))) / ((N(a)S)/N_a(3pl)) \ (N(3sg))S / N_a(L) \\
\hline
(N(3sg))S \wedge ((N(3sg))S)/N_a(3pl)) \\
\hline
\qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \wedge E_a \\
\qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad N(3sg)S
\end{array}$$

$$(46) \quad (\pi_1(x_{els} \ x_{escriu})) \\
\text{(write 3(sg))}$$

The first and second person clitics may also double with a realised prepositional phrase; the object of the co-occurring complement will be an agreeing reflexive such as *a mi (mateix)* ('myself'). This could be treated in a similar manner to the doubling of third person datives, though we do not pursue this here; that the non-third person cases carry always an emphatic force may suggest some distinction.

SUBJECT PRO-DROP

Main and embedded sentences of Catalan exhibit subject pro-drop, by which subjects are omitted; they are interpreted according to principles of agreement.

- (47) a. Canto.
'I sing'
- b. El Joan pensa que canto.
'John thinks that I sing'

In order to treat this we will employ explicit and semantically potent existential feature quantification, semantically interpreted by pairing and projection, like conjunction (see Morrill 1990a). The rules of inference for the existential type constructor are as follows.

$$\begin{array}{l}
(48) \quad \Gamma \Rightarrow \alpha : A[v \leftarrow F] \\
\hline
\Gamma \Rightarrow (F, \alpha) : \exists v A \\
\\
\Gamma \Rightarrow \gamma : \exists v A \quad \Delta(x : A[v \leftarrow (\pi_1 \gamma)]) \Rightarrow \beta : B \\
\hline
\Delta(\Gamma) \Rightarrow \beta[x \leftarrow (\pi_2 \gamma)] : B
\end{array}$$

Consider *que*. In order to allow its complement sentence to be without a subject it may have type $CP/(\exists a(N(a)S))$ where the embedded sentence is missing a nominal at its left edge with some agreement. The semantics is to take the existential type argument and apply the verbal meaning (second projection) to the agreement value (first projection): $(\lambda x((\pi_2 x) (\pi_1 x)))$.

$$\begin{array}{l}
(49) \quad \text{que} \quad \text{trobo} \quad \text{la Maria} \\
\hline
CP/(\exists a(N(a)S)) (N(1(sg))S)/N(_) N(3(sg)) \\
\hline
N(1(sg))S \\
\hline
\exists a(N(a)S) \\
\hline
CP
\end{array}$$

$$(50) \quad (x_{\text{que}} (1(sg), (x_{\text{trobo}} x_{\text{la Maria}}))) \\
((\text{find mary}) 1(sg))$$

To allow in addition an embedded sentence to be complete, a semantically potent disjunction will be used:

$$\begin{array}{l}
(51) \quad \Gamma \Rightarrow \alpha : A \quad \Gamma \Rightarrow \beta : B \\
\hline
\Gamma \Rightarrow (i\alpha) : A \vee B \quad \Gamma \Rightarrow (j\beta) : A \vee B \\
\hline
\Gamma(x : A) \Rightarrow \gamma : C \quad \Gamma(y : B) \Rightarrow \gamma_1 : C \\
\hline
\Gamma(w : A \vee B) \Rightarrow (w \rightarrow x. \gamma; y. \gamma_1) : C
\end{array}$$

The semantic interpretation is given in terms of a programming case statement: in the case that input is tagged *i*, substitute in the first branch; in the case *j*, in the second branch.

$$(52) \quad ((i\alpha) \rightarrow x. \gamma; y. \gamma_1) = \gamma[x \leftarrow \alpha] \\
((j\alpha) \rightarrow x. \gamma; y. \gamma_1) = \gamma_1[y \leftarrow \alpha]$$

$$(53) \quad \text{que} - (\lambda x(x \rightarrow y. ((\pi_2 y) (\pi_1 y)); z. z)) := CP/((\exists a(N(a)S)) \vee S)$$

Since subject pro-drop is also permitted in main clauses, where there is no embedding element to license the omission, it is proposed to admit a simple generalisation of the formalism:

$$(54) \quad \text{Main Clause Types}$$

Rather than a single distinguished sentence type, there is a parameter of grammar specifying a finite set of main clause type formulas.

Such a parameter (which in a categorial grammar of the kind envisaged here will be the only one other than the lexicon) would be employed in relation to main clause phenomena such as topicalisation, and V2. Thus in English, the main clause types will include NP·(S/NP), PP·(S/PP), etc. where · is the product. Then a main clause may consist of a topicalised maximal projection followed by a sentence lacking that constituent. An appropriate semantic term should be assigned to each main clause type formula showing how semantics in these pragmatically significant types is mapped into the logical value in the actual (truth-valued) sentence domain. In the case of Catalan main clause subject pro-drop there is the following:

$$(55) \quad \exists a(N(a)\backslash S) \Rightarrow ((\pi_2 x) (\pi_1 x))$$

For alternative presentations of subject pro drop, in terms of lexical lifting of nominals, as well as for unificational cliticisation, see Beaven (1990) and Sanfilippo (1990).

MEDIAL CLITICISATION

The types used so far command a rather too limited control over word order. They form a logic emphasising immediate adjacency in resource order, but are not suited to partial ordering constraints allowing limited permutation. In Morrill, *et al.* (1990) and Barry *et al.* (1991) a structural modality for permutation is proposed, drawing inspiration from the use of structural modalities in linear logic to govern resource transformations by structural rules. This machinery will be applied to allow non-peripheral cliticisation such as that in (56) which the earlier types are unable to generate.

$$(56) \quad \text{La dono a la Maria.}$$

‘I give it to Mary’

The categorization of *donar* is as illustrated in (57).

$$(57) \quad \text{dono - give} := ((N(1(\text{sg}))\backslash S)/N_a(_))/N(_)$$

Clitics are to now receive types such as (58) seeking a verb phrase missing a permutable nominal. The logic of permutation is presented in (59).

$$(58) \quad \text{la} - (\lambda x(x \text{ 3}(\text{sgf}))) := (N(a)\backslash S)/((N(a)\backslash S)/(HN(3(\text{sg}))))$$

$$(59) \quad \frac{\Gamma \Rightarrow \alpha : HA \quad H\Gamma \Rightarrow \alpha : A}{\Gamma \Rightarrow \alpha : A \quad H\Gamma \Rightarrow \alpha : HA} \begin{array}{c} \text{-----[HE]} \quad \text{-----[HI]} \\ \text{-----[HP], } A \text{ or } B \text{ is HD} \end{array}$$

$$\frac{\Gamma x : A \ y : B \ \Delta \Rightarrow \gamma : C}{\Gamma y : B \ x : A \ \Delta \Rightarrow \gamma : C}$$

The word order in (56) is now derived as follows:

- (65) $es \rightarrow et|us \rightarrow em|ens \rightarrow ho$
 $\rightarrow el|la|els|les \rightarrow en \rightarrow hi$
 $\rightarrow li$

The prescribed and Barcelona clitics differ from the València in function as well as order. In the course of presentation some indications will be made in the direction of capturing alternative orderings, but note that these remarks fall well short of addressing the alternative systems properly. There are many more subtleties.

We treat clitic order by means of a feature on sententials for clitisation class (cf. Baschung *et al.* 1987): a whole number which encodes increasing clitisation reactivity with increasing value. For space and clarity, clitic class feature structures are notated 1, 2, 3, ... and $\bar{1}$, $\bar{2}$, ... for 1, 2, 3, ... and not less than 1, not less than 2, ...; these stand for $s(0)$, $s(s(0))$, $s(s(s(0)))$, ... and $s(_)$, $s(s(_))$, ...

$$(66) \quad es \quad := \quad (N(3n)\backslash S(0))/((N(3n)\backslash S(1))/(HN(3(n))))$$

$$(67) \quad et \quad := \quad (N(a)\backslash S(1))/((N(a)\backslash S(2))/(HN(2(sg))))$$

$$us \quad := \quad (N(a)\backslash S(1))/((N(a)\backslash S(2))/(HN(2(pl))))$$

$$(68) \quad em \quad := \quad (N(a)\backslash S(2))/((N(a)\backslash S(3))/(HN(1(sg))))$$

$$ens \quad := \quad (N(a)\backslash S(2))/((N(a)\backslash S(3))/(HN(1(pl))))$$

$$(69) \quad li \quad := \quad (N(a)\backslash S(3))/((N(a)\backslash S(4))/(HN_a(3(sg))))$$

$$els \quad := \quad (N(a)\backslash S(3))/((N(a)\backslash S(4))/(HN_a(3(pl))))$$

$$(70) \quad el \quad := \quad (N(a)\backslash S(4))/((N(a)\backslash S(5))/(HN(3(sg))))$$

$$la \quad := \quad (N(a)\backslash S(4))/((N(a)\backslash S(5))/(HN(3(sg))))$$

$$els \quad := \quad (N(a)\backslash S(4))/((N(a)\backslash S(5))/(HN(3(pl))))$$

$$les \quad := \quad (N(a)\backslash S(4))/((N(a)\backslash S(5))/(HN(3(pl))))$$

$$(71) \quad en \quad := \quad (N(a)\backslash S(5))/((N(a)\backslash S(6))/(HN_d(3(_))))$$

$$(72) \quad hi \quad := \quad (N(a)\backslash S(6))/(N(a)\backslash S(7))$$

$$(73) \quad ho \quad := \quad (N(a)\backslash S(6))/((N(a)\backslash S(7))/(HCP))$$

For prescribed Catalan, the assignment to *li* in (69) becomes the following:

$$(74) \quad li \quad := \quad (N(a)\backslash S(3))/((N(a)\backslash S(5))/(HN_a(3(sg))))$$

Likewise, for the Barcelona dialect ordering the entries for *li* and *ho* should become:

$$(75) \quad li \quad := \quad (N(a)\backslash S(3))/((N(a)\backslash S(7))/(HN_a(3(sg))))$$

$$ho \quad := \quad (N(a)\backslash S(3))/((N(a)\backslash S(7))/(HCP))$$

And the clitic *els* now belongs to one class only. For a rationale of the Barcelona system see Gavarró (1990).

BLOCKING PREPOSITION STRANDING

Stranding of prepositions in Catalan is ungrammatical.

- (76) a. *Ens escrius a.
 'You(sg) write to us'
 b. *Ens parles de
 'You(sg) talk about us'

The types presented so far allow such overgeneration. Following Morrill (1990c) and Oehrle and Zhang (1989), island constraints will be captured by incorporating division duals to a non-associative product. Lambek (1961) presents the non-associative calculus (77); the division operators are written here as angles directed from domain to range.

$$(77) \quad \frac{\Gamma \Rightarrow \gamma : A < B \quad \Delta \Rightarrow \beta : B}{[\Gamma \Delta] \Rightarrow (\gamma \beta) : A} \text{[<E]} \qquad \frac{\Delta \Rightarrow \beta : B \quad \Gamma \Rightarrow \gamma : B < A}{[\Delta \Gamma] \Rightarrow (\gamma \beta) : A} \text{[>E]}$$

$$\frac{[\Gamma \ y : B] \Rightarrow \alpha : A}{\Gamma \Rightarrow (\lambda y \alpha) : A < B} \text{[<I]} \qquad \frac{[y : B \ \Gamma] \Rightarrow \alpha : A}{\Gamma \Rightarrow (\lambda y \alpha) : B > A} \text{[>I]}$$

While the associative calculus adheres to a list structure on resources, the non-associative one adheres to a binary tree structure. When the systems are mixed, sequents become partially bracketed ($n+2$ -ary trees) and non-associative inference is conditioned on the requisite bracketing (resource structure). The effect of the following non-associative assignments is to block clitics binding into the bracketed domain induced by prepositions: (79) is not a theorem, so a prepositional verb and a preposition do not form a member of a clitic's transitive verb argument type.

- (78) a - (λxx) := $N_a(a) < N(a)$
 de - (λxx) := $N_a(a) < N(a)$

(79) $(N(3(\text{sg})) \backslash S) / N_a(_)$ $N_a(a_1) < N(a_1) \Rightarrow (N(a) \backslash S) / N(3(\text{sg}))$

The use of non-associativity here is meant to carry a commitment to prosodic interpretation, rather than to be an ad hoc device. The idea is that partially bracketed structures designate constraints on prosodic constituency: necessary prosodic constituents. The treatment of prepositions as non-associative functors portrays prepositional phrases as obligatory prosodic constituents, and the effect of blocking stranding is derivative on this. In general the prediction is that domains which are necessarily prosodic constituents are islands to extraction, but this will not be the only constraint, and additional apparatus may allow more penetrative binders than the clitics.

CLAUSE-LOCALITY

Although a clitic does not necessarily bind an immediate argument of the verb to which it attaches, the position it binds must be local to its clause:

- (80) a. Penso que el Joan troba la Maria.
 b. *La penso que el Joan troba.

The grammar as it stands does not capture this clause-locality. The strategy of the previous section is not applicable since while preposition-stranding is to be blocked for all constructions, clauses are not ceilings for phenomena other than cliticisation such as relativisation. Following Morrill (1990b) we will read the relevant information on locality off the type system by including in it encoding of intensional semantic domains, in particular temporal domains. The unary temporally intensional type-constructor \langle is associated with semantic operations of intensionalisation and extensionalisation with respect to time indices; these satisfy 'down-up' cancellation:

$$(81) \quad (\downarrow(\uparrow\alpha)) = \alpha$$

In general the lexical semantics of words are now functions from time points to extensional denotations at those time points.

$$(82) \quad \text{troba - find} := \langle((N(3(\text{sg}))\backslash S)/N(_))$$

Taking proper names to be rigid designators, their lexical semantics will be constant functions, mapping time points into the same individual.

$$(83) \quad \text{el Joan - } (\uparrow\text{john}) := \langle N(3(\text{sg}))$$

The rules of inference are given in (84) and the derivation in (85) has the semantics (86).

$$(84) \quad \frac{\Gamma \Rightarrow \alpha : \langle A}{\Gamma \Rightarrow (\downarrow\alpha) : A} \text{[}\langle E \text{]} \qquad \frac{\langle \Gamma \Rightarrow \alpha : A}{\langle \Gamma \Rightarrow (\uparrow\alpha) : \langle A} \text{[}\langle I \text{]}$$

$$(85) \quad \begin{array}{c} \text{El Joan} \qquad \text{troba} \qquad \text{la Maria} \\ \hline \langle N(3(\text{sg})) \quad \langle((N(3(\text{sg}))\backslash S)/N(_)) \quad \langle N(3(\text{sg})) \\ \hline \text{-----}\langle E \text{-----}\langle E \\ (N(3(\text{sg}))\backslash S)/N(_) \quad N(3(\text{sg})) \\ \hline \text{-----}\langle E \text{-----}/E \\ N(3(\text{sg})) \qquad N(3(\text{sg}))\backslash S \\ \hline \text{-----}\backslash E \\ S \\ \text{---}\downarrow \\ \langle S \end{array}$$

- (86) a. $(\uparrow(((\downarrow x_{\text{troba}}) (\downarrow x_{\text{la Maria}})) (\downarrow x_{\text{El Joan}})))$
 b. $(\uparrow(((\downarrow \text{find}) \text{mary}) \text{john}))$

The [⟨I] in (85) is permitted since the derivation's components are of the form $\langle A$. Thus the sentence can adapt to the argument type of the intensional context creating element *que*.

(87) penso - **think** := <((N(1(sg))\S)/CP)
 que - ($\uparrow(\lambda xx)$) := <(CP/<S))

However, with the type assignment (88), a clitic will not bind into a temporal domain since its argument must take an element of type N(3(sg)) to form a verb phrase: an expression *penso que el Joan troba* requires an intensional type like <N(3(sg)).

(88) la - ($\uparrow(\lambda x(x\ 3(sg)))$) := <((N(a)\S)/((N(a)\S)/N(3(sg))))

CLITIC CLIMBING

Although control verbs require for semantic interpretation the intensions of their verb phrase complements across worlds, these are untensed and do not form temporal domains.

(89) puc := <((N(1(sg))\S)/VP)
 llegir := <(VP/N(_))

A word order like (90b) is prevented by assignments as in (89) since a prefix clitic is a functor over N(_)\S and not VP. The clitic climbing (90d) is expected, since there is no intervening tense inflected domain.

- (90) a. Puc llegir aquest llibre.
 'I am able to read this book'
 b. *Puc-lo llegir.
 c. Puc llegir-lo.
 d. El puc llegir.

Clitic climbing may extend across more than one verb:

- (91) a. Vull poder llegir aquest llibre.
 'I want to be able to read this book'
 b. El vull poder llegir

However, some verbs, such as *decidir*, cannot receive themselves clitics that have climbed (92d), and block climbing to superordinate verbs (93b).

- (92) a. Decideixo llegir aquest llibre.
 'I decide to read this book'
 b. *Decideixo-lo llegir.
 c. Decideixo llegir-lo.
 d. *El decideixo llegir.
 (93) a. Puc decidir llegir aquest llibre.
 'I am able to decide to read this book'
 b. *El puc decidir llegir

We propose here to classify those verbs allowing climbing as inheriting the cliticisation feature of their complement, and those blocking it as instantiating cliticisability to 0. Any blocking verb in a chain will prevent clitic climbing.

(94)	puc	:= <((N(1(sg))\S(i))/VP(i))
	poder	:= <(VP(i)/VP(i))
	decideixo	:= <((N(1(sg))\S(0))/VP(_))
	decidir	:= <(VP(0)/VP(_))

ACKNOWLEDGEMENTS

The research by the first author was partly supported by the Nederlandse Organisatie voor Wetenschappelijk Onderzoek, and that by the second was supported by the Ministerio de Educación y Ciencia del Estado Español. The origins of this paper lie with a visit of the first author to Clermont-Ferrant at the end of 1988.

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APPENDIX

The treatments of the various aspects of cliticisation in Catalan addressed in this paper are integrated in the following lexical assignments. The grammar has been implemented in the parser-theorem prover of Morrill (1990c).

el Joan	- (\uparrow john)	:= $\langle N(3(\text{sg})) \rangle$
la Maria	- (\uparrow mary)	:= $\langle N(3(\text{sg})) \rangle$
trobar	- find	:= $\langle (VP(\overline{7})/N(_)) \rangle$
trobo	- find	:= $\langle ((N(1(\text{sg})))\overline{S}(\overline{7}))/N(_) \rangle$
trobos	- find	:= $\langle ((N(2(\text{sg})))\overline{S}(\overline{7}))/N(_) \rangle$
troba	- find	:= $\langle ((N(3(\text{sg})))\overline{S}(\overline{7}))/N(_) \rangle$
trobem	- find	:= $\langle ((N(1(\text{pl})))\overline{S}(\overline{7}))/N(_) \rangle$
trobeu	- find	:= $\langle ((N(2(\text{pl})))\overline{S}(\overline{7}))/N(_) \rangle$
troben	- find	:= $\langle ((N(3(\text{pl})))\overline{S}(\overline{7}))/N(_) \rangle$
escriure - write		:= $\langle (VP(\overline{7})/N_a(_)) \rangle$
donar	- give	:= $\langle ((VP(\overline{7})/N_a(_))/N(_) \rangle$
pensar	- think	:= $\langle (VP(\overline{7})/CP) \rangle$
poder	- able	:= $\langle ((VP(i))/VP(i)) \rangle$
decidir	- decide	:= $\langle ((VP(0))/VP(_)) \rangle$
a	- ($\uparrow(\lambda xx)$)	:= $\langle (N_a(a)\langle N(a) \rangle) \rangle$
de	- ($\uparrow(\lambda xx)$)	:= $\langle (N_d(a)\langle N(a) \rangle) \rangle$
que	- ($\uparrow(\lambda x(\uparrow(\downarrow x) \rightarrow y. ((\pi_2 y) (\pi_1 y)); z. z))))$:= $\langle (CP/(\langle (\exists a(N(a)\overline{S}))\vee S)) \rangle$
jo	- ($\uparrow(\lambda x(x\ 1(\text{sg}))))$:= $\langle (S(i)/(N(1(\text{sg})))\overline{S}(i)) \rangle$
tu	- ($\uparrow(\lambda x(x\ 2(\text{sg}))))$:= $\langle (S(i)/(N(2(\text{sg})))\overline{S}(i)) \rangle$
ell	- ($\uparrow(\lambda x(x\ 3(\text{sgm}))))$:= $\langle (S(i)/(N(3(\text{sg})))\overline{S}(i)) \rangle$
ella	- ($\uparrow(\lambda x(x\ 3(\text{sgf}))))$:= $\langle (S(i)/(N(3(\text{sg})))\overline{S}(i)) \rangle$
nosaltres	- ($\uparrow(\lambda x(x\ 1(\text{pl}))))$:= $\langle (S(i)/(N(1(\text{pl})))\overline{S}(i)) \rangle$
vosaltres	- ($\uparrow(\lambda x(x\ 2(\text{pl}))))$:= $\langle (S(i)/(N(2(\text{pl})))\overline{S}(i)) \rangle$
ells	- ($\uparrow(\lambda x(x\ 3(\text{plm}))))$:= $\langle (S(i)/(N(3(\text{pl})))\overline{S}(i)) \rangle$

elles - ($\uparrow(\lambda x(x\ 3(plf)))$) := $\langle(S(i)/(N(3(pl))\S(i)))$

es - ($\uparrow(\lambda x(\lambda y((x\ y)\ y)))$) :=
 $\langle((N(3(n))\S(0))/(N(3(n))\S(1))/(HN(3(n))))$

et - ($\uparrow(\lambda x(x\ 2(sg)))$) :=
 $\langle((N(a)\S(1))/(N(a)\S(2))/(H(N(2(sg))\&N_a(2(sg))))$

us - ($\uparrow(\lambda x(x\ 2(pl)))$) :=
 $\langle((N(a)\S(1))/(N(a)\S(2))/(H(N(2(pl))\&N_a(2(pl))))$

em - ($\uparrow(\lambda x(x\ 1(sg)))$) :=
 $\langle((N(a)\S(2))/(N(a)\S(3))/(H(N(1(sg))\&N_a(1(sg))))$

ens - ($\uparrow(\lambda x(x\ 1(pl)))$) :=
 $\langle((N(a)\S(2))/(N(a)\S(3))/(H(N(1(pl))\&N_a(1(pl))))$

li - ($\uparrow(\lambda x((x\ 3(sg)),\ x))$) :=
 $\langle(((N(a)\S(3))\wedge(N(a)\S(3))/N_a(3(sg)))/(N(a)\S(4))/(HN_a(3(sg))))$

els - ($\uparrow(\lambda x((x\ 3(pl)),\ x))$) :=
 $\langle(((N(a)\S(3))\wedge(N(a)\S(3))/N_a(3(sg)))/(N(a)\S(4))/(HN_a(3(pl))))$

el - ($\uparrow(\lambda x(x\ 3(sg)))$) :=
 $\langle((N(a)\S(4))/(N(a)\S(5))/(HN(3(sg))))$

la - ($\uparrow(\lambda x(x\ 3(sg)))$) :=
 $\langle((N(a)\S(4))/(N(a)\S(5))/(HN(3(sg))))$

els - ($\uparrow(\lambda x(x\ 3(pl)))$) :=
 $\langle((N(a)\S(4))/(N(a)\S(5))/(HN(3(pl))))$

les - ($\uparrow(\lambda x(x\ 3(pl)))$) :=
 $\langle((N(a)\S(4))/(N(a)\S(5))/(HN(3(pl))))$

en - **- of-it** :=
 $\langle((N(a)\S(5))/(N(a)\S(6))/(HN_d(3(_))))$

hi - **- there** :=
 $\langle((N(a)\S(6))/(N(a)\S(7)))$

ho - **- so** :=
 $\langle((N(a)\S(3))/(N(a)\S(6))/(HCP)))$