

Some definitions:

12. Let  $A[1 \dots n]$  be an array with  $n$  different integers. Let  $ran(n)$  a randomized function that outputs an integer  $i$ ,  $1 \leq i \leq n$  under the uniform distribution. Let us consider the following algorithms

**Input:**  $A[1 \dots n]$   
**from**  $i = 1$  **to**  $n$  **do**  
     swap values  $A[i]$  and  $A[ran(n)]$

**Input:**  $A[1 \dots n]$   
**from**  $i = n$  **downto**  $1$  **do**  
     swap values  $A[i]$  and  $A[ran(i)]$

Do these algorithms output a uniform random permutation? Justify your answer.

13. Consider the following algorithm to generate an integer  $r \in \{1, \dots, n\}$ : We have  $n$  coins labelled  $m_1, \dots, m_n$ , where the probability that  $m_i = \text{head}$  is  $1/i$ . Toss the in order the coins  $m_n, m_{n-1}, \dots$  until getting the first head, if the fist head appears with coin  $m_i$ , the  $r = i$ . Prove that the previous algorithm yield an integer  $r$  with uniform distribution. i.e. the probability of getting any integer  $r$  is  $1/n$ .)
14. Consider the set  $S = \{1, \dots, n\}$ .
- We generate  $X \subseteq S$  as follows: A fair coin is flipped independently for each element of  $S$ , if the coin lands H, the element is added to  $X$ , otherwise it is not. Prove that the resulting set  $X$  is equally likely to be any one of the  $2^n$  possible subsets.
  - Suppose  $X, Y \subseteq S$  are chosen independently and u.a.r. from all  $2^n$  subsets from  $S$ . Compute  $\Pr[X \subseteq Y]$  and  $\Pr[X \cup Y = S]$
15. Let  $h(x) = x \bmod m$  be a hash function, where  $m = 2^p - 1$  for some prime number  $p$ . Let  $w$  be a character string corresponding to the representation in radix  $2^p$  of a natural number. Prove that, if  $w'$  is a string obtained by permuting the symbols in  $w$ ,  $h(w') = h(w)$ .
16. Consider the family  $\mathcal{H}$  of hash functions  $h : \{1, 2, 3, 4\} \rightarrow \{0, 1\}$  containing the three following functions
- $h_1(1) = 0, h_1(2) = 1, h_1(3) = 1, h_1(4) = 0$
  - $h_2(1) = 1, h_2(2) = 0, h_2(3) = 1, h_2(4) = 0$
  - $h_3(1) = 1, h_3(2) = 1, h_3(3) = 0, h_3(4) = 0$

Is  $\mathcal{H}$  universal? Justify your answer.

17. There are different operations that we wish to implement on sets of integers. Take into account that in a set there are no repeated elements. Let  $a[n]$  and  $b[n]$  be arrays holding the elements in two sets  $A$  and  $B$  with  $n$  elements. Provide an exact and a randomized algorithm for each of the following problems.

- (a) Given  $a$  and  $x \in \mathbb{N}$ , are there integers  $y, z \in A$  such that  $x = y + z$ ?
- (b) Given  $a$  and  $b$ , is  $A = B$ ?

Your algorithms should solve the problems in time  $\Theta(n \log n)$  and in expected time  $O(n)$ , respectively.

18. Lucas' theorem says the following: If we have an integer  $a$  such that:  $a^{n-1} \equiv 1 \pmod{n}$ , and, for every prime factor  $q$  of  $n - 1$ , it is not the case that  $a^{(n-1)/q} \equiv 1 \pmod{n}$ , then  $n$  is prime.

Can this result be used to show that Primality belongs to NP?

19. Consider the following problems:

- MODULAR FACTORIAL: Given  $N$  bits natural numbers  $x, y$  compute  $x! \pmod{y}$ .
  - SMALLEST PRIME DIVISOR: Given a  $N$  bit natural number  $x$ , compute the smallest prime divisor of  $x$ .
  - FACTORING: Given a  $N$  bit natural number  $x$ , compute the factorization of  $x$  as product of primes.
- (a) Prove that  $y$  is prime if and only if, for each integer  $x < y$ , we have that  $\text{mcd}(x!, y) = 1$ .
  - (b) Show that if MODULAR FACTORIAL can be solved in polynomial time, then SMALLEST PRIME DIVISOR and FACTORING could be solved in polynomial time.