

The *CatLog3* Technical Manual

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Chapter 1

Preamble

Logical categorial grammar is a version of categorial grammar in which the syntactic calculus is a universal categorial logic, i.e. the syntactic component is invariant from language to language, and a language is described entirely by its lexicon, and a grammar is purely lexical. *CatLog* is a series of programs for automatic logical categorial grammar semantic parsing (that is, the computer translation of human language into logic) developed by the present author; hitherto the programs have been written in Prolog. *CatLog3* is a system for such semantic parsing based on the particular program *CatLog* version l2.pl.

For execution, the engine *CatLog3* requires a logical categorial lexicon, and some linguistic examples. Running *CatLog3* on a lexicon and some examples supports textual listing of a lexicon and generation of linguistic analyses in the Prolog console: the linguistic expressions, associated sequent calculus derivations/syntactic structures, and logical forms. Such execution also creates files containing L^AT_EX code for the lexicon, and L^AT_EX code for outputs for the examples resulting from *CatLog3* linguistic analysis (the linguistic expressions, associated sequent calculus derivations/syntactic structures, and logical forms).

The structure of this technical manual is as follows. In Chapter 2 the parser/theorem-prover *CatLog3* is introduced. Chapter 3 illustrates execution outputting to the Prolog console. Chapter 4 lists the standard, predefined, linguistic examples of the file `examples.pl` of *CatLog3*, and Chapter 5 lists the standard, predefined, English lexicon of the file `lexicon.pl` of *CatLog3*. In Chapter 6 the LaTeX output is exemplified. The file `l2.pl` annexing this technical manual contains the Prolog source code of version l2.pl of *CatLog*, the program implementing parsing/theorem-proving, on the basis of what we call ‘hedge’-sequent calculus, for the fragment of categorial type logic of *CatLog3* (Morrill 2017[11]); the files `example.pl` and `lexicon.pl` also annexing this technical manual contain the Prolog source code of the English linguistic examples and lexicon respectively.

Chapter 2

Introduction

The *CatLog* programs work by reducing grammar to logic: an expression α is of the grammatical type A if and only if an associated logical statement $Lex(\alpha) \Rightarrow A$, where $Lex(\alpha)$ is a lexical selection for α , is a theorem of a universal categorial type calculus. This type calculus is an intuitionistic sublinear logic. It is *sublinear* in that *theoremhood is not fully preserved under structural rules such as contraction or permutation* (word multiplicity and word order matter in grammar). It is *intuitionistic* in that its proofs are *constructive*, and the reading of a derived expression is a function of the meanings of its lexical expressions and of the constructive content of the proof by which it is deemed to be grammatical. I.e. the design entails *compositionality*, the principle which is generally formulated as stating that the meaning of a sentence is a function of the meanings of its words and of their mode of composition; this principle is attributed to the founder of modern logic G. Frege, although he never stated it as such; it's dual would appear to be *contextuality*: the principle according to which the meaning of a sentence is also a function of its 'context' (idea communicated to me by H. Hendriks).

This overall architectural design is purely lexical: a grammar is just a lexicon (dictionary) which is a set of ternary lexical items (signs) of the form:

lex(Prosodic form, Type, Logical form).

The Prosodic form is a string of words which may be partially bracketed; the brackets represent prosodic or syntactical domains such as intonational phrases or extraction islands. The Prosodic form is the phonosyntactical representation of the lexical item. The Type is a formula of the categorial type logic; the type represents the grammatical category of the lexical item. The Logical form is a closed term of higher-order logic representing the logical semantics of the lexical item. The language model defined by a set of lexical signs is the set of ternary signs derivable from/generated by this lexicon Lex according to the logic of the categorial grammatical types:

$$\{(\alpha, A, \phi) \mid \vdash Lex(\alpha) \Rightarrow A : \phi\}$$

As well as being called logical categorial grammar, grammar on this design has been called, for example, type logical categorial grammar, type logical grammar, or logical grammar; see, for example, Moortgat 1988[6], 1997[8]; Morrill 1994[14], 2011[15]; Carpenter 1997[1]; Jäger 2005[3]; Oehrle 2011[17]; or Moot and Retoré 2012[9]).

The task of semantic parsing is that of computing, for a given phonosyntactical expression α , and grammatical category A , the set of semantic readings (logical forms) of the expression α in type A according to a grammar.

The finite reading property (van Benthem 1991[18]) is the property according to which in each type an expression is associated with at most a finite number of logical forms. This seems to be an empirical property of human language: there appear to be no infinitely ambiguous expressions or, at least, no expressions which are infinitely ambiguous within a single grammatical category.

Clearly, in the absence of the finite reading property explicit semantic parsing is not computable since printing the members of an infinite set can never terminate.¹

CatLog3 runs under SWI-Prolog version 7.6.4. There are three Prolog source files: the categorial calculus parser/theorem-proving engine `l2.pl`, the dictionary `lexicon.pl` of the substantial English fragment, and the small English corpus `examples.pl`. *CatLog3* computes, given an index of the identifiers of linguistic items in `examples.pl`, the expressions' associated logical forms according to the categorial dictionary of `lexicon.pl` and the categorial calculus of `l2.pl`. As well as pretty printing dictionaries and analyses in the Prolog console, *CatLog3* produces dictionary and analysis output files in the format of L^AT_EX typesetting code (based on Donald Knuth's T_EX).

The complete list of files used in *CatLog3* is:

- `l2.pl`: the Prolog source coding the parser/theorem-prover for the categorial type logic of *CatLog3*;
- `lexicon.pl`: the Prolog source coding the dictionary of the logical categorial grammar of English of *CatLog3*, in which each lexical entry is of the form:

`lex(Prosodic form, Type, Logical form)`

where *Prosodic form* is a string of words which may be partially bracketed, *type* is a formula of the categorial type logic of *CatLog3*, and *Logical form* is a closed term of higher-order logic;

- `examples.pl`: the Prolog source coding the linguistic corpus of English examples of *CatLog3*, in which each linguistic example is of the form:

`str(Identifier, Prosodic form, Type)`

where *Identifier* is a Prolog term identifying the example, *Prosodic form* is the expression to be parsed, in the form of a string of words which may be partially bracketed, and *type* is the target category of the parsing, in the form of a formula of the categorial type logic of *CatLog3*;

- `lex.tex`: a derived L^AT_EX file coding a lexicon, created and generated by loading `l2.pl` and a dictionary/lexicon (such as `lexicon.pl`) into Prolog, and executing the query:

?- pplexlatex.

¹In the categorial logic of *CatLog3* there are two sources of possible non-compliance with the finite reading property and of possible consequent undecidability: the subexponentials $!, ?$, inspired by the exponentials of Girard (1987[2]), (which interact with the bracket modalities $\langle \rangle, []^{-1}$ of Morrill 1992[10] and Moortgat 1996[7]), and the continuous product unit I and discontinuous product unit J of Lambek (1988[5]) and Morrill, Valentín, and Fadda (2011[13]) respectively. With respect to the units I and J , consider unit assignments in the lexicon (for 'null elements', for example), a lexical entry such as $0 : S/S : f$ in Lambek grammar, or $1 : S\uparrow S : f$ in displacement grammar. Both can make every sentence generated by the grammar infinitely ambiguous since these null functors can reapply to a sentence any number of times, changing it semantically but not phonosyntactically. And with respect to subexponentials $!$, Kanovich, Kuznetsov and Scedrov 2017[4] show undecidability of a bracket-conditioned contraction subexponential in the context of the Lambek calculus with bracket modalities.

Let us say that an occurrence of a connective in a lexical type is positive if it will (eventually) show up as the principal connective of an succedent type in sequent proof search, and that it is negative if it will (eventually) show up as the principal connective of an antecedent formula. In relation to possible non-computability we propose the following. Regarding the second source of infinite ambiguity above, to guarantee the finite reading property it is sufficient to require that in *CatLog3* lexical types are *bracket non-negative* (Morrill and Valentín 2015[12]) meaning that there must be no negative bracket modality $\langle \rangle$, nor any positive anti-bracket modality $[]^{-1}$, within any occurrence of a negative contraction subexponential $!$. And regarding the first source above, we propose, as in Morrill (2018[16]), that *CatLog3* lexical types be *unit non-negative* meaning that they have no continuous or discontinuous unit in negative position. So far as we are aware, in *CatLog3*, bracket and unit non-negativity, together with the assumption that lexical types have no negative expansion subexponentials $?$, are sufficient to ensure the finite reading property and for the fragment of categorial logic.

- `der.tex`: a derived L^AT_EX file coding the output of *CatLog3* parsing, created and generated by loading into Prolog `l2.pl`, a dictionary/lexicon (such as `lexicon.pl`), and some linguistic examples (such as `examples.pl`), and by executing a query:

```
?- tt(Index).
```

where *Index* is a Prolog term unifying with identifiers of the linguistic examples wanted;

- `out.tex`: a L^AT_EX file which includes the files `lex.tex` and `der.tex`, and which when typeset produces the file `out.pdf` containing the final formatted dictionary and analyses resulting from loading into Prolog `l2.pl`, a dictionary/lexicon, and some linguistic examples, and executing the queries:

```
?- pplexlatex.
```

```
?- tt(Index).
```

- `prooftree.sty`: Paul Taylor's macro package for formatting proof trees in L^AT_EX, used by `out.tex`.

Listouts of the SWI-Prolog source files `l2.pl`, `lexicon.pl` and `examples.pl` of *CatLog3* are given in the files `l2.pl`, `lexicon.pl`, and `examples.pl`, annexed to this technical manual.

Chapter 3

Outputting to the Prolog console

In *CatLog3* the program *CatLog* version l2 implementing parsing/theorem-proving for the categorial type calculus is in the file `l2.pl`, and a logical categorial grammar of English is represented by the dictionary in the file `lexicon.pl`. In the file `examples.pl` there is a corpus of linguistic examples which are expressions (prosodic forms) that are partially bracketed strings of English words.

3.1 Outputting of a dictionary to the Prolog console

In *CatLog3*, once `l2.pl` and `lexicon.pl` are loaded into SWI-Prolog, the query:

```
?- pplex.
```

pretty prints in the Prolog console the categorial dictionary in `lexicon.pl`. This is as follows:


```

everyone: nLAF((sfciAght(g))inSF): "LARB((vperson B) -> (A B))
face: nLCNs(n): face
fell: nL(Ea-<Na>SF): "LA(Past (vfall A))
filled: nL(<=>Eght(s(g))\SF)/EaNa: "LALB(Past ((vfile A) B))
finds: nL(<=>Eght(s(g))\SF)/EaNa: "LALB(Pres ((vfind A) B))
fish: nLCNs(n): fish
for: iL(PPfor/EaNa): LAA
form: nL(CNs(n)&nt(s(m))): "vform. (gen vform)
fortunately: nLAF(vsfInSF): fortunately
friends: nL(Cmp/PPof): Friends
from: nL((AaAF(<=>Na>SF)\(<=>Na>SF)\&nt(Cm)&nt(Cm))\EBNB): "LA((vfromadv A), (vfromadv A))
gave: nL((<=>EaNa>SF)/EBNB-PPto): "LALB(Past (((vgive p1ZA) p1IA) B))
gave: nL(<=>Eght(s(g))\SF)/EaNa_wfthe_cold_shoulder): "LALB(Past ((vshun A) B))
gave: nL((<=>EaNa>SF)/EaNa): "LALB(CPast (((vgive A) B) C))
girl: nLCNs(f): girl
gives: nL(<=>Eght(s(g))\SF)/EaNa_wfthe_cold_shoulder): "LALB(Pres ((vshun A) B))
God: iLnt(s(m)): God
good: nLnt(Cm)&nt(Cm): good
has: nL(<=>Eght(s(g))\SF)/EaNa: "LALB(Pres ((vhave A) B))
he: iL[-I]_Iag((ILSg Ica iLnt(s(m)))/(<=>nt(s(m))\Sg)): LAA
heaven: nLCNs(m): heaven
her: iLAgAa((<=>Na>Sg)ciLnt(s(f))in(iL(<=>Na>Sg) Ica iLnt(s(f)))): LAA
himself: iLAF((<=>nt(s(m))\SF)\ciLnt(s(m))\in(<=>nt(s(m))\SF)): "LALB((A B)
horse: nLCNs(n): horse
in: nL(AaAF(<=>Na>SF)\(<=>Na>SF)/EaNa): "LALB(C((vin A) (B C))
in: nL(AF(SfmsD)/EaNa): in
is: iL(<=>Eght(s(g))\SF)/EaNa_Eg((CNg/CNg)|+|(CNg/CNg)-I)): LALB(Pres (A->C.[B = C]. D.(D LE[E = B] B)))
it: iLW(i): 0
it: iLAFaA((<=>Na>SF)\ciLnt(s(m))\in(iL(<=>Na>SF) Ica iLnt(s(m)))): LAA
it: iL[-I]_Iaf((ILSf Ica iLnt(s(m)))/(<=>nt(s(m))\SF)): LAA
jogs: nL(<=>Eght(s(g))\SF): "LA(Pres (vjog A))
John: iLnt(s(m)): j
laughs: nL(<=>Eght(s(g))\SF): "LA(Pres (vlaugh A))
left: nL(<=>Eght(s(g))\SF): "LA(Pres (vleave A))
let: nL(Sim/Sb): let
light: nL(CNs(m)&nt(s(m))): "vlight. (gen vlight)
likes: nL(<=>Eght(s(g))\SF)/EaNa: "LALB(Pres ((vlike A) B))
logic: nL(NT(s(n))&Cns(n)): "(gen vlogic), vlogic
London: iLnt(s(m)): l
loses: nL(<=>Eght(s(g))\SF)/EaNa: "LALB(Pres ((vlose A) B))
love: nL(<=>EaNa(Sb)/EaNa): "LALB((vlove A) B)
loved: nLAaAb((<=>Na>S->ciNB)dp((<=>Na>S->ciNB)imAg(CNg/CNg))): "vlove. LALB(C((B C) & ED((A C) D)))
loves: nL(<=>Eght(s(g))\SF)/EaNa: "LALB(Pres ((vlove A) B))
man: nLCNs(m): man
mary: iLnt(s(f)): m
more: nL((<=>EaNa>SF)/EaNa): "LALB(Past ((vmeet A) B))
mountain: nLCNs(m): mountain
mountain: nLCNs(m): mountain
necessarily: iL(CA/mLsA): nec
of: nL((nt(Cm)&nt(Cm))/iLENB)&(Pfor/EaNa)): "(vof, LMA)
or: iLAAf((ealLSF[-I]-I]_ISF)/iLSD: [mapphin.0.or]
or: iLAAf((ealLSF(<=>Na>SF)[-I]-I]_I(<=>Na>SF))/iL(<=>Na>SF): [mapphin.[app.s.0].or]
or: iLAAf((ealLSF(<=>Eght(s(g))\SF))[-I]-I]_ISF(<=>Eght(s(g))\SF))/iL(SF(<=>Eght(s(g))\SF)): [mapphin.[app.s.0].or]
or: iLAAf((ealLSF(<=>Na>S->ciNB)dp((<=>Na>S->ciNB)imAg(CNg/CNg))[-I]-I]_I((<=>Na>SF)/EBNB))/iL((<=>Na>SF)/EBNB)): [mapphin.[app.s.[app.s.0]].or]
painting: nL(CNs(n)/Pfor): LA((vof A) vpainting)
paper: nLCNs(n): paper
park: nLCNs(n): park
past: nLAaAF((<=>Na>SF)\(<=>Na>SF)/EBNB): "LALB(C((vpast A) (B C))
perseverance: nL(NT(s(n))&Cns(n)): "(gen vperseverance), vperseverance)
peter: iLnt(s(m)): p
phonetics: nL(NT(s(n))&Cns(n)): "(gen vphonetics), vphonetics)
praises: nL(<=>Eght(s(g))\SF)/EaNa: "LALB(Pres ((vpraise A) B))
raced: nL(<=>EaNa>SF): "LA(Past (vrace A))
raced: nLAaAb((<=>Na>S->ciNB)dp((<=>Na>S->ciNB)imAg(CNg/CNg))): "vrace2. LALB(C((B C) & ED((A C) D)))
rains: nL(<=>W(i)-osf): "Cpres vitrains)
reading: nL((<=>EaNa>Spp)/EaNa): "LALB((vread A) B)
robin: iLght(s(g)): f
said: nL((<=>EaNa>SF)/Sib): "LALB(Past ((vsay A) B))
saw: nL((<=>EaNa>SF)/EaNa-CPht): "LALB(Past ((A->C.(vsee C) D.(vsee D) B))

```

```

seeks: nl((=>Egnt(s(g))SE)/nLAaaf((<<na\SE)/EBhb)\(na\SE)): "LALB(vtries ^((va vfind) B) B)
sees: nl((=>Egnt(s(g))SE)/EaHa): "LALB(Pres ((vsee A) B))
sent: nl((=>EaHa\SE)/EBhb^PPro): "LALB(Past ((vsent pizza) p1IA) B)
sent: nl((=>EaHa\SE)/EaHa): "LALB(Past ((vsend A) B) C)
she: :IL[]-IAG(ILSg Ica iLrt(s(f)))/(<>rt(s(f))Sg): LAA
sings: nl((=>Egnt(s(g))SE): "LA(Pres (vsing A))
slept: nl((=>Egnt(s(g))SE): "LA(Past (vsleep A))
slowly: nLAaafGt((<<na\SE)\(<<nlNa\SE)): "LALB(vslowly ^ (va VB))
sneezed: nl((=>Egnt(s(g))SE): "LA(Past (vsneeze A))
sold: nl((=>EaHa\SE)/EBhb^PPro): "LALB(Past ((vsell pizza) p1IA) B)
someone: nLAaaf(SfciILight(g))inSD: "LAEB(vpersion B) & (A B)
Spirit: iLCns(0): Spirit
studies: nl((=>Egnt(s(g))SE)/EaHa): "LALB(Pres (vstudy A) B)
such-that: iLAn(CMn/Cm)/Sf Ica iLrt(n)): "LALB(C(B C) & (A C))
suzy: iLrt(s(f)): s
talks: nl((=>Egnt(s(g))SE): "LA(Pres (vwalk A))
tall: nLAg(Cng/Cng): tall
teetotal: nAn(CMn/Cm): "LALB(A B) & (vteetotal B)
temmiliondollars: iLrt(s(n)): temmiliondollars
than: iLCPHany/ALSf: LAA
that: iLAn(1-1]-1(CMn/Cm)/iL((<>rt(n)|&ueiLrt(n))\SE): "LALB(C(B C) & (A C))
the: iLAn(0n)/Cm): iora
the-cold+shoulder: iLW[the.cold.shoulder]:
there: iLW[there]:
thinks: nl((=>Egnt(s(g))SE)/CPthat+HLSE): "LALB(Pres ((vthink A) B))
to: iL(PPro/EaHa)|&An((<>nm\SD)/(<<nm\SB)): LAA
today: nLAaaf((<<na\SE)\(<<na\SE)): "LALB(vtoday (A B))
tries: nl((=>Egnt(s(g))SE)/nL((=>Egnt(s(g))SE)): "LALB(vtries ^ (va B) B)
unicorn: nLCns(n): unicorn
up: iLW[up]:
upon: nl((Abaf((<<nb\SE)\(<<nb\SE)WAg(Cng/Cng)/EaHa): "LA((vuponadv A), (vuponadv A))
void: nLAg(Cng/Cng): void
walk: nl((=>EaHa-Egnt(s(g))SE): "LA(Pres (vwalk A))
walk: nl((=>EaHa\SB): "LA(vwalk A)
walks: nl((=>Egnt(s(g))SE): "LA(Pres (vwalk A))
was: iL((=>Egnt(s(g))SE)/EaHa+EG(Cng/Cng)+I(Cng/Cng))-I)): "LALB(Past (A->C.[B = C]; D.<(D LE[E = B]) B)))
was: nl((=>W[there]-oSE)/EaHa): "LA(Past (obe A))
waters: nLCmp(0): waters
which: iLAnM(0n)clne(n))in(1-1]-1(CMn/Cm)/iL((<>rt(n)|&ueiLrt(n))\SE)): "LALB(CID(C D) & (B (A D)))
will: iLAA((<<na\SE)/(<<na\SB)): "LALB(Fut (A B))
without: nLAg(Cng/Cng)/EaHa): "LALB(C(B C) & -(with A) C)
without: iLAAf(1-1((<<na\SE)\(<<na\SE)))/(<<na\Spap)): "LALB(C(B C) & -(A C))
woman: nLCns(f): woman
yesterday: nLAaaf((<<na\SE)\(<<na\SE)): "LALB(vyesterday (A B))
true.

```

3.2 Outputting of analyses to the Prolog console

The query

```
?- tt(N).
```

parses all the examples the index of which unifies with the Prolog term N , and times the task. For example, a session with `?- tt(1)` is as follows, with `examples.pl`, `lexicon.pl`, and `l2.pl` as in the annexes to this technical manual already consulted.

```
MacBook-Air-de-Glyn:CatLog3 morrill$ swipl
Welcome to SWI-Prolog (threaded, 64 bits, version 7.6.4)
SWI-Prolog comes with ABSOLUTELY NO WARRANTY. This is free software.
Please run ?- license. for legal details.
```

```
For online help and background, visit http://www.swi-prolog.org
For built-in help, use ?- help(Topic). or ?- apropos(Word).
```

```
Cannot read termcap database;
using dumb terminal settings.
?- [l2].
This is CatLog3. Copyright 2018 (C) Glyn V. Morrill.
CatLog3 comes with ABSOLUTELY NO WARRANTY. This is free software.
true.
```

```
?- [examples].
true.
```

```
?- [lexicon].
true.
```

```
?- tt(1).
```

```
(1) [john]+walks Sf
```

```
[iLNt(s(m)): j], nL(<>EgNt(s(g))\Sf): ^LA(Pres (vwalk A)) => Sf
```

```
[iLNt(s(m))], nL(<>EgNt(s(g))\Sf) => Sf [LiL]
  [iLNt(s(m))], <>EgNt(s(g))\Sf => Sf [\iL]
    [iLNt(s(m))] => <>EgNt(s(g)) [<>R]
      iLNt(s(m)) => EgNt(s(g)) [ER]
        iLNt(s(m)) => Nt(s(m)) [iLiL]
          Nt(s(m)) => Nt(s(m))
            Sf => Sf
```

```
(Pres (vwalk j))
% 4,942 inferences, 0.003 CPU in 0.004 seconds (86% CPU, 1609772 Lips)
true
```

We see first the prosodic form (bracketed string) and the type in which it is to be analysed. Then there is the associated logic statement a proof of which would show grammaticality, including the

lexical semantics of the basic expressions. Then there is such a proof. Finally there is the logical translation that results from all this.

Chapter 4

examples.pl

```

str(dwp('7-7'), [b(john), walks], s(F)).
str(dwp('7-16'), [b(every, man)], talks, s(F)).
str(dwp('7-19'), [b(the, fish)], walks, s(F)).
str(dwp('7-32'), [b(every, man)], b([b(every, man)], b([walks, or, talks]))], s(F)).
str(dwp('7-34'), [b(b([every, man], walks, and, b([she]), talks))], s(F)).
str(dwp('7-39'), [b(b([a, woman]), walks, and, b([she]), talks))], s(F)).
str(dwp('7-43, 45'), [b(john)], believes, that, b([a, fish]), walks, s(F)).
str(dwp('7-48, 49, 52'), [b(every, man)], believes, that, b([a, fish]), walks, s(F)).
str(dwp('7-57'), [b(every, fish, such, that, b(it), walks)], talks, s(F)).
str(dwp('7-69, 62'), [b(john)], seeks, a, unicorn, s(F)).
str(dwp('7-73'), [b(john)], is, bill, s(F)).
str(dwp('7-76'), [b(john)], is, a, man, s(F)).
str(dwp('7-83'), [necessarily, b(john)], walks, s(F)).
str(dwp('7-86'), [b(john)], walks, slowly, s(F)).
str(dwp('7-91'), [b(john)], tries, to, walk, s(F)).
%str(dwp('7-94'), [b(john)], tries, to, b([catch, a, fish, and, eat, it]))], s(F)).
str(dwp('7-98'), [b(john)], finds, a, unicorn, s(F)).
str(dwp('7-105'), [b(every, man, such, that, b(the), loves, a, woman]), loses, her], s(F)).
str(dwp('7-110'), [b(john)], walks, in, a, park, s(F)).
str(dwp('7-116, 118'), [b(every, man)], doesn't, walk], s(F)).

str(1, [b(john)], walks, s(F)).
str(2, [b(john)], loves, mary], s(F)).
str(3, [b(john)], thinks, b(mary)], walks, s(F)).
str(4, [b(mary)], believes, b(john)], thinks, b(mary)], walks, s(F)).
str(5, [b(mary)], buys, john, coffee], s(F)).
str(6, [man, loved], cn(s(m))).
str(7, [b(john)], is, loved, s(F)).
str(8, [b(john)], is, loved, by, mary], s(F)).
str(9, [b(the, man)], walks, s(F)).
str(10, [b(the, man)], loves, the, woman], s(F)).
str(11, [man, b([b(that, walks]))], cn(s(m))).
str(12, [b(the, man, b([b(that, walks]))], talks, s(F)).
str(13, [b(the, man, b([b(that, loves, mary]))], walks, s(F)).
str(14, [b(the, man, b([b(that, b(mary), loves]))], walks, s(F)).
str(15, [b(the, man, b([b(that, b(john), thinks, b(mary), loves]))], walks, s(F)).
str(16, [b(the, horse)], raced, past, the, barn], s(F)).
str(17, [b(the, horse, raced, past, the, barn)], fell], s(F)).
str(18, [b(john)], walks, from, edinburgh], s(F)).
str(19, [b(the, man, from, edinburgh)], walks], s(F)).

str(20, [b(hond)], is, '007'], s(F)).
str(21, [b(hond)], is, teetotal], s(F)).
str(22, [b(hond)], is, b([b('007', and, teetotal))], s(F)).
str(23, [b(hond)], is, b([b(teetotal, and, '007'))], s(F)).
str(24, [b(hond)], is, b([b(tall, and, '007')], s(F)).

str(tdc(43), [b(mary)], gave, the, man, the, cold, shoulder], s(F)).
str(tdc(44), [b(john)], gave, john, the, cold, shoulder], s(F)).
str(tdc(47), [b(john)], gave, every, book, to, mary], s(F)).
str(tdc(50), [b(mary)], thinks, b([someone]), left], s(F)).
str(tdc(53), [b(everyone)], loves, someone], s(F)).
str(tdc(56), [dog, b([b(that, b(mary), saw, today))], cn(s(n))].
str(tdc(58a), [b(john)], slept, before, b(mary), did], s(F)).
%str(tdc(58b), [john, slept, and, mary, did, too], s(F)).
str(tdc(69), [mountain, b([b(the, painting, of, which, by, cezanne, b(john), sold, for, tenmilliondollars))], cn(s(n))].
str(56, [thesis, the, height, of, the, lettering, on, the, first, line, of, the, second, page, of, the, third, chapter, of, which, is, '0.5cm'], cn(n)).
str(tdc(67), [b(john, b([b(who, jogs))], sneezed, s(F)).
str(tdc(70a), [b(john)], fortunately, has, perseverance], s(F)).
str(tdc(70c), [b(john)], fortunately, has, perseverance], s(F)).
str(tdc(70d), [b(john)], has, fortunately, perseverance], s(F)).
%str(tdc(73), [b(b([b(john), studies, logic, and, b(charies)], phonetics))], s(F)).
str(tdc(86a), [b(john)], ate, more, donuts, than, b(mary), bought, bagsels], s(F)).
str(tdc(86b), [b(john)], bought, himself, coffee, s(F)).
%str(65, [dorothy, bet, the, straw, man, half, of, himself, that, she, would, reach, emerald, city, first], s).

str(w(1), [b(it), rains], s(F)).

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str(w(2), [b(mary)], gives, peter, the, cold, shoulder], s(f)).
str(w(3a), [b(mary)], calls, peter, up], s(f)).
str(w(3b), [b(mary)], calls, up, peter], s(f)).
%str(w(4a), [b(b(peter)], sees, mary, and, b(robin)], clark]]], s(f)).
%str(w(4b), [b(b(peter)], calls, up, mary, and, b(robin)], clark]]], s(f)).
%str(w(4c), [b(b(peter)], calls, mary, up, and, b(robin)], clark]]], s(f)).
%str(w(4d), [b(b(peter)], gives, mary, the, cold, shoulder, and, b(robin)], clark]]], s(f)).

str(eacrel(6), [man, b(b(that, walks))], cn(s(m))].
str(eacrel(7), [man, b(b(that, b(mary), likes))], cn(s(m))].
str(eacrel(8), [b(john)], likes, the, man], s(f)).
str(eacrel(9), [b(mary)], thinks, that, b(john)], likes, the, man], s(f)).
str(eacrel(10), [b(suzy)], believes, that, b(mary)], likes, the, man], s(f)).
str(eacrel(11), [man, b(b(that, b(john)], likes))], cn(s(m))].
str(eacrel(12), [man, b(b(that, b(mary)], thinks, that, b(john)], likes, the, man], s(f)).
str(eacrel(13), [man, b(b(that, b(suzy)], believes, that, b(mary)], thinks, that, b(john)], likes, the, man], s(f)).
str(eacrel(14), [man, b(b(that, b(suzy)], talks, and, b(bill)], s(f)).
str(eacrel(15), [b(b(mary)], talks, and, b(bill)], s(f)).
str(eacrel(16), [b(b(john)], talks, and, b(bill)], s(f)).
str(eacrel(17), [b(b(suzy)], talks, and, b(bill)], s(f)).
str(eacrel(18), [b(b(john)], talks, and, b(bill)], s(f)).
str(eacrel(19), [man, that, b(the, friends, of)], walk], cn(s(m))].
str(eacrel(20), [man, that, b(the, friends, of, admire))], cn(s(m))].
str(eacrel(21), [paper, b(b(that, b(john)], filed, without, reading))], cn(s(n))].
str(eacrel(22), [paper, b(b(that, the, editor, of, filed, without, reading))], cn(s(n))].
%str(eacrel(23), [b(b(that, the, editor, of, filed, without, reading))], cn(s(n))].

str(eacrnd(14), [b(b(b(john)], praises, mary, and, b(john)], laughs))], s(f)).
str(eacrnd(15), [b(john)], b(b(praises, mary, and, laughs))], s(f)).
str(eacrnd(16), [b(john)], b(b(likes, and, will, love))], london], s(f)).
str(eacrnd(17), [b(john)], b(b(gave, or, sent))], mary, the, book], s(f)).
str(eacrnd(18), [b(b(b(john)], or, b(mary))], sings], s(f)).
str(eacrnd(19), [b(john)], loves, b(b(mary, and, himself))], s(f)).
str(eacrnd(20), [b(b(b(john)], likes, and, b(mary), loves))], london], s(f)).
str(eacrnd(21), [b(john)], b(b(gave, the, book, and, sent, the, cd))], to, mary], s(f)).
str(eacrnd(22), [b(john)], gave, b(b(the, book, to, mary, and, the, cd, to, suzy))], s(f)).
str(eacrnd(23), [b(john)], saw, b(b(mary, today, and, bill, yesterday))], s(f)).
%str(eacrnd(24), [man, b(b(that, b(b(john)], saw, yesterday, and, met, today))], cn(s(m))].
str(eacrnd(25), [man, b(b(that, b(b(john)], saw, yesterday, and, met, today))], cn(s(m))].
str(eacrnd(26), [b(b(b(john)], walks, b(mary)], talks, and, b(bill)], s(f)).
str(eacrnd(27), [b(john)], b(b(walks, talks, and, sings))], s(f)).
str(eacrnd(28), [b(john)], b(b(praises, likes, and, will, love))], london], s(f)).

str(atb(1), [man, b(b(that, b(b(b(john)], likes, and, b(mary), loves)))]], cn(s(m))].
str(atb(2), [man, b(b(that, b(mary)], b(b(likes, and, praises)))]], cn(s(m))].
str(gen(1), [in, the, beginning, b('God')], created, b(b(the, heaven, and, the,
earth))], s(f)).
%str(gen(2('a(1))), ['And', b(b(b(the, earth)], was, b(b(without, form, and, void))], and, b(darkness)], was, upon, the, face, of, the, deep)]], s(f)).
str(gen(2('a(1))), [b(the, earth)], was, b(b(without, form, and, void))], s(f)).
str(gen(2('a(1))), [b(darkness)], was, upon, the, face, of, the, deep], s(f)).
str(gen(2('b(1))), ['And', b(the, 'spirit', of, 'God')], moved, upon,
the, face, of, the, waters], s(f)).
str(gen(3), ['And', b(b(b('God')], said, let, b(tthere)], be, light,
and, b(tthere)], was, light)]], s(f)).
%str(gen(4('a(1))), ['And', b('God')], saw, b(b(the, light, and, that, b(tthere)], was, good)]], s(f)).
%str(gen(4('aa(1))), ['And', b('God')], saw, b(b(the, light,
and, that, b(the, light)], was, good)]], s(f)).
% str(gen(4('b(1))), [b('God')], divided, the, light, from, the, darkness], s(f)).

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% str(Gen(one('a'))), ['And', b([b([b(['God']), called, the, light, day, and, the, darkness, he, called, night]])], s(f)).
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Chapter 5

lexicon.pl

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lex({'007'}, il g iu n(c(s(g))), '007').
lex({a}, il g iu (f iu (s(f) ci il n(c(s(g)))) in s(f)/cn(s(g))), [lmd, X, [lmd, Y, [ext, Z, [and, [app, X, Z], [app, Y, Z]]]]]).
lex({admirer}, il (br(c(a ie n(c)) - (g ie n(c(s(g)))) bs s(f)/(a ie n(a))), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [dh, admirer], X], Y]]]]]).
lex({'And'}, il f iu (s(f)/s(f)), [lmd, X, X].
lex({and}, il a iu f iu(see il A bs ab ab A)/il A), [mapPhin, [app, s, Ø], and] :-
  A = s(f)/ue n(a);
  A = br n(a)bs s(f);
  A = f iu(see il A bs ab A)/il A), [mapPhin, [app, s, Ø], and] :-
  A = s(f)/a ie n(a);
  A = s(f) ci((br n(a) bs s(f)/rio w(q)/n(b)))licn w(q).
lex({and}, il f iu a iu(see il A bs ab ab A)/il A), [mapPhin, [app, s, Ø], and] :-
  A = (br n(a) bs s(f)/b ie n(b));
  A = (br n(a) bs s(f)/b ie n(b))bs(br n(a) bs s(f));
  A = f iu a iu(il A bs ab ab A)/il A), [lmd, X, [lmd, Y, [lmd, Z, [and, [app, Y, Z], W], [app, [app, X, Z], W]]]]];
  A = (br n(a) bs s(f)/(c ie n(c)+cp(b)))bs(br n(a)bs s(f));
  A = (br n(a)bs s(f)/pp(b));
  A = (br n(a)bs s(f)/(c ie n(c)'pp(b)))bs(br n(a) bs s(f));
  A = (br n(a)bs s(f)/ue n(b)).
lex({ate}, nl (br a ie n(a) bs s(f)/a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [dh, eat], X], Y]]]]]).
lex({bagels}, nl (n(c(p(n)))&n(p(n))), [up, [pair, [app, gen, [dh, bagels]], [dh, bagels]]]).
lex({barn}, nl cn(s(n)), barn).
lex({be}, nl (br w(there) iu s(b))/a ie n(a)), [up, [lmd, X, [app, [dh, be], X]]]).
lex({before}, il(a iu f iu (br n(a) bs s(f))bs (br n(a) bs s(f))/s(f)), [lmd, X, [lmd, Y, [lmd, Z, [app, [app, before, X], [app, Y, Z]]]]]).
lex({beginning}, nl cn(s(n)), beginning).
lex({believes}, nl (br g ie n(c(s(g))) bs s(f)/(cp(that) iad nl s(f))), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [dh, believe], X], Y]]]]]).
lex({bill}, il n(c(s(m))), b).
lex({bond}, il n(c(s(m))), b).
lex({book}, nl cn(s(n)), book).
lex({bought}, nl (br a ie n(a) bs s(f)/(a ie n(a)))*, [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dh, buy], [fst, X], [snd, X], Y]]]]]).
lex({bought}, nl (br a ie n(a) bs s(f)/a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dh, buy], X], Y]]]]]).
lex({buy}, il(a iu (br n(a) bs s(f))bs (br n(a) bs s(f))n(a)), [lmd, X, [lmd, Y, [lmd, Z, [and, [eq, Z, X], [app, Y, Z]]]]]).
lex({buy}, nl (n iu cn(n) bs cn(n))/a ie n(a)), [up, [lmd, X, [lmd, Y, [app, [dh, buy], X], Y]]]).
lex({buys}, nl (br g ie n(c(s(g))) bs s(f)/(a ie n(a)))*, [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [dh, buy], [fst, X], [snd, X], Y]]]]]).
lex({call}, nl (br g ie n(c(s(g))) bs s(f)/a ie n(a))lip n(a))iad(n(a)rip w(lup))), [up, [lmd, X, [lmd, Y, [app, [dh, phone], X], Y]]]).
lex({catch}, nl (br a ie n(a) bs s(b))/a ie n(a)), [up, [lmd, X, [lmd, Y, [app, [dh, catch], X], Y]]]).
lex({cezanne}, il n(c(s(m))), c).
lex({cd}, nl cn(s(n)), cd).
lex({charles}, il n(c(s(m))), c).
lex({clark}, il g iu n(c(s(g))), c).
lex({coffee}, nl (n(c(s(n)))&n(s(n))), [up, [pair, [app, gen, [dh, coffee]], [dh, coffee]]]).
lex({created}, nl (br a ie n(a) bs s(f)/a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dh, create], X], Y]]]]]).
lex({darkness}, nl cn(s(n))&n(c(s(n))), [up, [pair, [dh, darkness], [app, gen, [dh, darkness]]]).
lex({deep}, nl cn(s(n)), deep).
lex({did}, il a iu g iu b iu h iu (((br n(a) bs s(g)) ci (br n(b) bs s(b)))/(c ie br n(c) bs s(f))bs((br n(a) bs s(g))ci(n( ) bs s(b))))), [lmd, X, [lmd, Y, [app, [app, X, Y], Y]]]).
lex({did, too}), ((br n(c) bs s(c)) ci (br n(c) bs s(c))bs((br n(c) bs s(c))ci(n( ) bs s(c))))), [lmd, X, [lmd, Y, [app, [app, X, Y], Y]]]).
lex({doesnt}, il g iu a iu (s(g)ci((br n(a)bs s(f)/(br n(a)bs s(b))))in s(g)), [lmd, X, [lmd, Y, [lmd, Z, [app, Y, Z]]]]]).
lex({dog}, nl cn(s(n)), dog).
lex({donuts}, nl (n(c(p(n)))&n(p(n))), [up, [pair, [app, gen, [dh, donuts]], [dh, donuts]]]).
lex({earth}, nl cn(s(n)), earth).
lex({eat}, nl (br a ie n(a) bs s(b))/a ie n(a)), [up, [lmd, X, [lmd, Y, [app, [app, [dh, eat], X], Y]]]]]).
lex({edinburgh}, il n(c(s(n))), e).
lex({editor}, nl (g iu cn(s(g))/pp(o)), editor).
lex({everyone}, il g iu (f iu (s(f) ci n(c(s(g)))) in s(f)/cn(s(g))), [lmd, X, [lmd, Y, [all, Z, [imply, [app, X, Z], [app, Y, Z]]]]]).
lex({face}, nl cn(s(n)), face).
lex({fall}, nl (a ie br n(a)bs s(f)), [up, [lmd, X, [app, 'Past', [app, [dh, fall], X]]]]]).
lex({filed}, nl (br g ie n(c(s(g))) bs s(f)/a ie n(c)), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [dh, file], X], Y]]]]]).
lex({finds}, nl (br g ie n(c(s(g))) bs s(f)/a ie n(c)), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [dh, find], X], Y]]]]]).
lex({fish}, nl cn(s(n)), fish).
lex({fish}, il(pp(top)/a ie n(a)), [lmd, X, X].
lex({form}, nl cn(s(n))&n(c(s(n))), [up, [pair, [dh, form], [app, gen, [dh, form]]]).
lex({fortunately}, nl f iu (sp s(f)in s(f)), fortunately).

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lex({{friends}}, nl (cn(p)/pp(ot)), friends).
lex({{from}}, nl ((a iu f iu (br n(a)bs s(f))&(n iu (cn(m)bs cn(n)))/(b ie n(b)))). [up, [lmd, X, [pair, [app, [dn, fromadv], X], [app, [dn, fromadv], X]]]]].
%lex({is}, il(g ie n(t(s(g)))bs s(f))/(g ie (cn(g)bs cn(g))))). [up, [lmd, X, [lmd, Y, [app, [app, X, [lmd, Z, [eq, Z, Y]]], Y]]]].
lex({have}), nl (br a ie n(a)bs s(f))/(b ie n(b)/pp(to)), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dn, give], [snd, X]], [fst, X]], [fst, X]], Y]]]]].
lex({have}), nl ((br g ie n(t(s(g)))bs s(f))/(a ie n(a) rip w(the, cold, shoulder))), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dn, shun], X], Y, Z]]]]]].
lex({have}), nl ((br g ie n(a)bs s(f))/(a ie n(a)))/(a ie n(a)), [up, [lmd, X, [lmd, Y, [lmd, Z, [app, 'Past', [app, [app, [dn, give], X], Y, Z]]]]]].
%lex({give}, i, the, cold, shoulder), (n(c)bs s(f) ci n(c), shun).
lex({give}), nl (cn(s(f)), girl).
lex({give}), nl (br g ie n(t(s(g)))bs s(f))/(a ie n(a) rip w(the, cold, shoulder))), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [app, [dn, shun], X], Y]]]]]].
lex({'God'}, il n(t(s(m))), 'God').
lex({good}), nl n iu (cn(n)/cn(n)), good).
lex({has}), nl (br g ie n(t(s(g)))bs s(f))/(a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [app, [dn, have], X], Y]]]]]].
lex({he}), il ab g iu ((l s(g) lca il n(t(s(m)))/(br n(t(s(m)))bs s(g))))). [lmd, X, X]].
lex({heaven}), nl (cn(s(n)), heaven).
lex({here}), il g iu a ((br n(a)bs s(g) ci il n(t(s(f))) in (il (br n(a)bs s(g) lca il n(t(s(f))))), [lmd, X, X])).
lex({himself}), il f iu ((br n(t(s(m)))in (il (br n(t(s(m)))bs s(f))ci n(t(s(m)))in (il (br n(a)bs s(g) lca il n(t(s(f))))), [lmd, X, X])).
lex({horse}), nl (cn(s(m)), horse).
lex({in}), nl (f iu (f iu (br n(a)bs s(f))/(a ie n(a)), [up, [lmd, X, [lmd, Y, [lmd, Z, [app, [app, [dn, in], X], [app, Y, Z]]]]]]]].
lex({is}), il ((br g ie n(t(s(g)))bs s(f))/(a ie n(a))-g ie ((cn(g)/cn(g))iad(cn(g)bs cn(g))>)), [lmd, X, [lmd, Y, [eq, Y, Z], W, [app, [app, W, [lmd, U, [eq, U, Y]]], Y]]]]]].
%lex({it}, il f iu ((s(f) ci il n(t(s(m))) in (il s(f) lca il n(t(s(m)))). [lmd, X, X])).
lex({it}), il f iu a ((br n(a)bs s(f) ci il n(t(s(m))) in (il (br n(a)bs s(f) lca il n(t(s(m)))). [lmd, X, X])).
lex({it}), il ab f iu ((l s(f) lca il n(t(s(m)))/(br n(t(s(m)))bs s(f))), [lmd, X, X])).
lex({jog}), nl (br g ie n(t(s(g)))bs s(f)), [up, [lmd, X, [app, 'Pres', [app, [dn, jog], X]]]]]].
lex({joke}), il n(t(s(m))), j).
lex({laughs}), nl (br g ie n(t(s(g)))bs s(f)), [up, [lmd, X, [app, 'Pres', [app, [dn, laugh], X]]]]]].
lex({left}), nl (br g ie n(t(s(g)))bs s(f)), [up, [lmd, X, [app, 'Pres', [app, [dn, leave], X]]]]]].
lex({let}), nl (s(im)/s(b)), let).
lex({light}), nl (cn(s(m))&n(t(s(n))), [up, [pair, [dn, light], [app, gen, [dn, light]]]]].
lex({likes}), nl (br g ie n(t(s(g)))bs s(f))/(a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [app, [dn, like], X], Y]]]]]].
lex({logic}), nl (cn(s(m))&n(s(n))), [up, [pair, [app, gen, [dn, logic]], [dn, logic]]].
lex({London}), il n(t(s(m))), L).
lex({loses}), nl (br g ie n(t(s(g)))bs s(f))/(a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [dn, lose], X], Y]]]]]].
lex({love}), nl (br a ie n(a)bs s(b)/a ie n(a)), [up, [lmd, X, [lmd, Y, [app, [app, [dn, love], X], Y]]]]]].
lex({loved}), nl (a iu (b iu ((br n(a)bs s(-)ci n(b))&(br n(a)bs s(-)ci n(b))in(g iu (cn(g)bs cn(g)))). [up, [pair, [dn, love], [lmd, X, [lmd, Y, [lmd, Z, [and, [app, Y, Z], [fst, W, [app, [app, X, Z], W]]]]]]]]]].
lex({loves}), nl (br g ie n(t(s(g)))bs s(f))/(a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [dn, love], X], Y]]]]]].
lex({man}), nl (cn(s(m)), man).
lex({marry}), il n(t(s(f))), m).
lex({meet}), nl (br a ie n(a)bs s(f))/(a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dn, meet], X], Y]]]]]].
lex({more}), il h g iu f iu ((s(f) ci il n(t(s(f))&(p(ot)tham) ci il (t)))). [lmd, X, [lmd, Y, [more-tham, [card, [lmd, Z, [app, X, [lmd, P, [lmd, R, [and, [app, P, Z], [app, R, Z]]]]]]]], [card, [lmd, Z], [dn, [app, Y, [lmd, P], [lmd, Q], [and, [app, P], Z]], [app, Q, Z]]]]]]]]]]. :- Q-(s(b) ci n(t(s(g))) in s(b))/cn(g).
lex({mountain}), nl (cn(s(n)), mountain).
lex({moved}), nl (br a ie n(a)bs s(f)), [up, [lmd, X, [app, 'Past', [app, [dn, move], X]]]]]].
lex({necessarily}), il(s(f)/nl s(f)), "nec").
lex({of}), nl ((n iu (cn(n)bs cn(n))/IL (b ie n(b))&(pp(ot) a ie n(a)))). [up, [pair, [dn, of], [lmd, X, X]]]].
lex({of}), il f iu ((ee il s(f)bs ab ab s(f))/IL s(f)), [mapphin, 0, of]).
lex({or}), il a iu f iu ((ee il A bs ab ab A)/IL A), [mapphin, [app, s, 0], of]) :- A = br n(a)bs s(f).
lex({or}), il f iu ((ee il A bs ab ab A)/IL A), [mapphin, [app, s, 0], of]) :-
A = s(f)/(br g ie n(t(s(g)))bs s(f)).
lex({or}), il a iu f iu ((ee il A bs ab ab A)/IL A), [mapphin, [app, s, [app, s, 0]], of]) :-
A = ((br n(a)bs s(f))/(b ie n(b)))/(b ie n(b)).
lex({painting}), nl (cn(s(n))/pp(ot)), [up, [lmd, X, [app, [app, [dn, of], X], [dn, painting]]]]].
lex({paper}), nl (cn(s(n)), paper).
lex({park}), nl (cn(s(n)), park).
lex({past}), nl a iu f iu ((br n(a)bs s(f)bs (br n(a)bs s(f)))/(b ie n(b))). [up, [lmd, X, [lmd, Y, [lmd, Z, [app, [app, [dn, past], X], [app, Y, Z]]]]]]]].
lex({perseverance}), nl (n(t(s(m))&n(s(n))), [up, [pair, [app, gen, [dn, perseverance]], [dn, perseverance]]]].
lex({person}), il n(t(s(m))), p).
lex({phonetics}), nl (n(t(s(n))&n(s(n))), [up, [pair, [app, gen, [dn, phonetics]], [dn, phonetics]]]].
lex({praises}), nl (br g ie n(t(s(g)))bs s(f))/(a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [app, [dn, praise], X], Y]]]]]].
%lex({raced}), ((cn(g)bs cn(g))/(a(b)bs n(b)bs s(-)))/(a(b)bs n(b)bs s(-)).
%pair, [lmd, X, [lmd, Y, [lmd, Z, [and, [app, Y, Z], [fst, W, [app, [app, X, Z], W]]]]]]]]. raceZ]].
lex({raced}), nl (a iu (b iu ((br n(a)bs s(-)ci n(b))&(br n(a)bs s(-)ci n(b))in(g iu (cn(g)bs cn(g)))). [up, [pair, [dn, raceZ], [lmd, X, [lmd, Y, [lmd, Z, [and, [app, Y, Z], [fst, W, [app, [app, X, Z], W]]]]]]]].
lex({rains}), nl (br w(t(t)) iu s(f)), [up, [app, 'Pres', [dn, it rains]]]].

```

lex((reading), nl (br a ie n(a) bs s(psp))/a ie n(a), [up, [lmd, X, [app, [app, [dn, read], X], Y]]]).
lex((robin), il g iu n((s(g))), r).
lex((said), nl (br a ie n(a) bs s(f))/s(lm), [up, [lmd, X, [lmd, Y, [app, 'Past',
[app, [app, [dn, say], X], Y]]]]]).
lex((saw), nl (br a ie n(a) bs s(f))/(a ie n(a)-cp(thatt)), [up, [lmd, X, [lmd, Y,
[app, 'Past', [app, [case, X, X], [app, [dn, see], X], X], [app, [dn, see], X], Y]]]).
lex((seeks), nl (br g ie n((s(g))) bs s(f))/nl a iu f iu ((n(a) bs s(f))/b ie n(b)) bs n(a) bs s(f))), [up, [app, [app, [dn, tries], [up, [app, [app, [dn, X], [dn, find], Y]], Y]]]).
lex((sees), nl (br g ie n((s(g))) bs s(f))/a ie n(a), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [app, [dn, see], X], Y]]]]]).
lex((sent), nl (br a ie n(a) bs s(f))/(b ie n(b)*pp(to)), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dn, sent], [snd, X]], [fst, X]], [fst, X]], Y]]]).
lex((sent), nl ((br (a ie n(a)) bs s(f))/(a ie n(a)))/(a ie n(a)), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dn, send], X], Y], Z]]]]]).
lex((she), il g iu ((l s(g) lca il n((s(f))))/(br n((s(f)))) bs s(g)), [lmd, X, X]).
lex((sings), il br g ie n((s(g))) bs s(f), [up, [lmd, X, [app, 'Pres', [app, [dn, sing], X]]]).
lex((sleep), nl (br g ie n((s(g))) bs s(f), [up, [lmd, X, [app, 'Past', [app, [dn, sleep], X]]]).
lex((slowly), nl a iu f iu(nl (br n(a) bs s(f)) bs (br nl n(a) bs s(f))), [up, [lmd, X, [lmd, Y, [app, [dn, slowly], [up, [app, [dn, X], [dn, Y]]]]]]]).
lex((sneezed), nl (br g ie n((s(g))) bs s(f), [up, [lmd, X, [app, [dn, sneeze], X]]]).
lex((soad), nl (br a ie n(a) bs s(f))/(b ie n(b)*pp(for)), [up, [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dn, seal], [snd, X]], [fst, X]], [fst, X]], Y]]]).
lex((someone), nl f iu ((s(f) ci il g iu n((s(g))) in s(f), [up, [lmd, X, [sst, Y, [and, [app, [dn, person], Y], [app, X, Y]]]]]).
lex((spirit), il cn(s(m)), 'spirit').
lex((studies), nl (br g ie n((s(g))) bs s(f))/a ie n(a), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [app, [dn, study], X], Y]]]).
lex((such, that), il n iu ((cn(n) bs cn(n))/(s(f) lca il n((n))), [lmd, X, [lmd, Y, [lmd, Z, [and, [app, Y, Z], [app, X, Z]]]]]).
lex((suz), il n((s(f))), s).
lex((talks), nl (br g ie n((s(g))) bs s(f), [up, [lmd, X, [app, 'Pres', [app, [dn, talk], X]]]).
lex((tall), nl (g iu (cn(g)/cn(g))), tall).
lex((teetotal), nl n iu (cn(n)/cn(n)), [up, [lmd, X, [lmd, Y, [and, [app, X, Y], [app, [dn, teetotal], Y]]]]]).
lex((temilliondollars), il n((s(n))), temilliondollars).
lex((than), il(gp(than)/l s(f)), [lmd, X, X]).
lex((that), il(n iu (ab ab (cn(n)bs cn(n))/(br n((n))iac ue il n((n))bs s(f))), [lmd, X, [lmd, Y, [lmd, Z, [and, [app, Y, Z], [app, X, Z]]]]]).
%lex((that), (cn(g)bs cn(g))/(s(f)/n((g))), [lmd, X, [lmd, Y, [lmd, Z, [and, [app, Y, Z], [app, X, Z]]]]]).
lex((the), il(n iu (n((n)/cn(n))), iota).
lex((the, cold, shoulder), il w((the, cold, shoulder)), 0).
lex((there), il w((there)), 0).
lex((thinks), nl (br g ie n((s(g))) bs s(f))/(cp(thatt)nl s(f)), [up, [lmd, X, [lmd, Y, [app, 'Pres', [app, [dn, think], X], Y]]]).
lex((to), il(gp(to)/a ie n(a))iac(n iu (br n(n) bs s(c)))/(br n(n) bs s(b))), [lmd, X, X]).
lex((today), nl a iu f iu (br n(a)bs s(f))/(br n(a)bs s(f)), [up, [lmd, X, [lmd, Y, [app, [dn, today], [app, X, Y]]]]]).
lex((tries), nl (br g ie n((s(g))) bs s(f))/nl (br g ie n((s(g))) bs s(f))), [up, [lmd, X, [lmd, Y, [app, [app, [dn, tries], [up, [app, [dn, X], [dn, Y]]]]]).
lex((uncorn), nl cn(s(n)), uncorn).
lex((up), il w((up)), 0).
lex((upon), nl (b iu f iu (br n(b) bs s(f))bs(br n(b) bs s(f)))&g iu (cn(g) bs cn(g))/a ie n(a), [up, [lmd, X, [pair, [app, [dn, uponadv], X], [app, [dn, uponadv], X]]]).
lex((void), nl g iu (cn(g)/cn(g)), void).
lex((walk), nl (br ((a ie n(a)) - (g ie n((s(g)))) bs s(f), [up, [lmd, X, [app, 'Pres', [app, [dn, walk], X]]]).
lex((walks), nl (br a ie n(a)bs s(b)), [up, [lmd, X, [app, [dn, walk], X]]]).
lex((walks), nl (br g ie n((s(g))) bs s(f), [up, [lmd, X, [app, 'Pres', [app, [dn, walk], X]]]).
lex((was), il((br g ie n((s(g))) bs s(f))/(a ie n(a)))/(a ie n(a))-(g ie ((cn(g)/cn(g))iac(cn(g) bs cn(g))-i)), [lmd, X, [lmd, Y, [app, 'Past', [app, [app, [dn, be], X]]]).
lex((was), nl (br w((there)) iu s(f))/a ie n(a), [up, [lmd, X, [app, 'Past', [app, [dn, be], X]]]).
lex((waters), nl cn(p(n)), waters).
lex((which), il(n iu m iu (cn(n))ci n((m)))in (ab ab (cn(n)bs cn(m))/(br n((n))iac ue il n((n))bs s(f))), [lmd, W, [lmd, Y, [lmd, Z, [and, [app, Y, Z], [app, X, [app, W, Z]]]]]).
lex((who), il (n iu n iu (ab ab (cn(n)bs(s(b) ci n((n))in s(b)))/(br n((n))iac ue il n((n))bs s(f))), [lmd, X, [lmd, Y, [lmd, Z, [and, [app, X, Y], [app, X, Y]]]).
lex((will), il a iu (br n(a)bs s(f))/(br n(a)bs s(b)), [lmd, X, [lmd, Y, [app, 'Fut', [app, X, Y]]]).
lex((without), nl (g iu (cn(g) bs cn(g))/a ie n(a)), [up, [lmd, X, [lmd, Y, [lmd, Z, [and, [app, Y, Z], [not, [app, [app, [dn, with], X], Z]]]]]).
lex((without), il a iu f iu(ab(br n(a)bs s(f))bs(br n(a)bs s(f)))/(br n(a) bs s(sp))), [lmd, X, [lmd, Y, [lmd, Z, [and, [app, Y, Z], [not, [app, X, Z]]]]]).
lex((woman), nl cn(s(f)), woman).
lex((yesterday), nl a iu f iu (br n(a)bs s(f))bs(br n(a)bs s(f)), [up, [lmd, X, [lmd, Y, [app, [dn, yesterday], [app, X, Y]]]]]).
onecoord(A) :- member(A, []).

```


Chapter 6

L^AT_EX output

Consider consulting `l2.pl`, `examples.pl` and `lexicon.pl` and then executing the queries:

```
?- pplexlatex, tt(-).
```

Invoking L^AT_EX on `out.tex` yields here:

```

007 :  $\mathbb{N}gNt(s(g)) : 007$ 
a :  $\mathbb{N}SgYf((Sf \mathbb{N}Ht(s(g)))Sf)CNs(g) : \lambda A \lambda B \text{EC}[(A \ C) \wedge (B \ C)]$ 
admire :  $\square((\exists \text{EaNa} \neg \exists gNt(s(g)))Sf) \text{EaNa} : \sim \lambda \lambda A \lambda B \text{Pres}(\text{"admire } A \ B)$ 
And :  $\mathbb{N}Yf(SfSf) : \lambda A \ A$ 
and :  $\mathbb{N}Yf(\exists Sf \sqcap \Pi^{-1} S f) \mathbb{N}Sf : (\Phi^{s+} \ 0 \ \text{and})$ 
and :  $\mathbb{N}Yaf(\exists ((\exists \text{Na} \ S f) \sqcap \Pi^{-1} \Pi^{-1} (\exists \text{Na} \ S f))) \mathbb{N}(\exists \text{Na} \ S f) : (\Phi^{s+} \ (s \ 0) \ \text{and})$ 
and :  $\mathbb{N}Yaf(\exists ((\exists \text{Na} \ S f) \sqcap \Pi^{-1} \Pi^{-1} (Sf \setminus \text{Na})) \mathbb{N}(Sf \setminus \text{Na})) : (\Phi^{s+} \ (s \ 0) \ \text{and})$ 
and :  $\mathbb{N}Yf(\exists (Sf \ \text{EaNa}) \sqcap \Pi^{-1} \Pi^{-1} (Sf \setminus \text{EaNa})) \mathbb{N}(Sf \setminus \text{EaNa}) : (\Phi^{s+} \ (s \ 0) \ \text{and})$ 
and :  $\mathbb{N}WtaVtaVYf(\mathbb{N}(Sf \setminus ((\exists \text{Na} \ S f) \rightarrow Wta) \sqcap Wta)) \sqcap \mathbb{N}Wta : \lambda \lambda A \lambda B \lambda C[(B \ C) \wedge (A \ C)]$ 
and :  $\mathbb{N}Yta((\exists \text{Na} \ S f) \ \text{EaNa}) \sqcap \Pi^{-1} \Pi^{-1} ((\exists \text{Na} \ S f) \ \text{EaNa}) : (\Phi^{s+} \ (s \ 0) \ \text{and})$ 
and :  $\mathbb{N}Yta((\exists \text{Na} \ S f) \ \text{EaNa}) \sqcap \Pi^{-1} \Pi^{-1} ((\exists \text{Na} \ S f) \ \text{EaNa}) : (\Phi^{s+} \ (s \ 0) \ \text{and})$ 
and :  $\mathbb{N}Yta(\mathbb{N}((\exists \text{Na} \ S f) \ \text{EaNa}) \setminus ((\exists \text{Na} \ S f) \ \text{EaNa})) \setminus \mathbb{N}((\exists \text{Na} \ S f) \ \text{EaNa}) : (\Phi^{s+} \ (s \ 0) \ \text{and})$ 
and :  $\mathbb{N}Yta(\mathbb{N}((\exists \text{Na} \ S f) \ \text{EaNa}) \setminus ((\exists \text{Na} \ S f) \ \text{EaNa})) \setminus \mathbb{N}((\exists \text{Na} \ S f) \ \text{EaNa}) : (\Phi^{s+} \ (s \ 0) \ \text{and})$ 
and :  $\mathbb{N}Yta((\exists \text{Na} \ S f) \ \text{EaNa}) \setminus \mathbb{N}((\exists \text{Na} \ S f) \ \text{EaNa}) \setminus \mathbb{N}((\exists \text{Na} \ S f) \ \text{EaNa}) : (\Phi^{s+} \ (s \ 0) \ \text{and})$ 
and :  $\mathbb{N}Yta((\exists \text{Na} \ S f) \ \text{EaNa}) \setminus \mathbb{N}((\exists \text{Na} \ S f) \ \text{EaNa}) \setminus \mathbb{N}((\exists \text{Na} \ S f) \ \text{EaNa}) : (\Phi^{s+} \ (s \ 0) \ \text{and})$ 
ate :  $\square((\exists \text{EaNa} \ S f) \ \text{EaNa}) : \sim \lambda \lambda A \lambda B \text{Past}(\text{"eat } A \ B)$ 
bagels :  $\square(\text{Nt}(\text{pr}) \setminus \text{CNpr}(\text{n})) : \sim (\text{gen } \text{"bagels"}, \text{bagels})$ 
barn :  $\square \text{CNs}(\text{n}) : \text{barn}$ 
be :  $\square((\exists \text{W}[\text{there}] \rightarrow S \text{b}) \ \text{EaNa}) : \sim \lambda A(\text{"he } A)$ 
before :  $\mathbb{N}Yaf((\exists \text{Na} \ S f) \setminus (\exists \text{Na} \ S f)) \setminus S f : \lambda \lambda A \lambda B \lambda C(\text{before } A \ B \ C)$ 
beginnings :  $\square \text{CNs}(\text{n}) : \text{beginning}$ 
believes :  $\square((\exists gNt(s(g)))Sf) \setminus (\text{CPHatt} \sqcup \text{Sf}) : \sim \lambda \lambda A \lambda B \text{Pres}(\text{"believe } A \ B)$ 
bill :  $\mathbb{N}Ht(s(m)) : b$ 
bond :  $\mathbb{N}Ht(s(m)) : b$ 
book :  $\square \text{CNs}(\text{n}) : \text{book}$ 
bought :  $\square((\exists \text{EaNa} \ S f) \ \text{EaNa}) \setminus \mathbb{N}(\text{EaNa}) : \sim \lambda \lambda A \lambda B \text{Past}(\text{"buy } \pi_1 A \ \pi_2 A \ B)$ 
bought :  $\square((\exists \text{EaNa} \ S f) \ \text{EaNa}) \setminus \mathbb{N}(\text{EaNa}) : \sim \lambda \lambda A \lambda B \text{Past}(\text{"buy } A \ B)$ 
by :  $\mathbb{N}Na((\exists \text{Na} \ S \rightarrow) \setminus (\exists \text{Na} \ S \rightarrow)) \setminus \text{Na} : \lambda \lambda A \lambda B \lambda C[C = A] \wedge (B \ C)$ 
by :  $\square(Yt(\text{CNa} \ \text{CNn}) \ \text{EaNa}) \setminus \mathbb{N}(\text{EaNa}) : \sim \lambda \lambda A \lambda B(\text{"by } A \ B)$ 
buys :  $\square((\exists gNt(s(g)))Sf) \setminus (\text{EaNa} \ \text{EaNa}) : \sim \lambda \lambda A \lambda B \text{Pres}(\text{"buy } \pi_1 A \ \pi_2 A \ B)$ 
calls :  $\square((\exists gNt(s(g)))Sf) \setminus \text{EaNa} \setminus (\text{W} \sqcup \text{P}) \setminus \text{Na} : \sim \lambda \lambda A \lambda B \text{Pres}(\text{"call } A \ B)$ 
catch :  $\square((\exists \text{EaNa} \ S \text{b}) \ \text{EaNa}) \setminus \mathbb{N}(\text{EaNa}) : \sim \lambda \lambda A \lambda B(\text{"catch } A \ B)$ 
cezanne :  $\mathbb{N}Ht(s(m)) : c$ 
cd :  $\square \text{CNs}(\text{n}) : \text{cd}$ 
charles :  $\mathbb{N}Ht(s(m)) : c$ 
clark :  $\mathbb{N}gNt(s(g)) : c$ 
coffee :  $\square(\text{Nt}(\text{s}(\text{n})) \setminus \text{CNs}(\text{n})) : \sim (\text{gen } \text{"coffee"}, \text{"coffee"})$ 
created :  $\square((\exists \text{EaNa} \ S f) \ \text{EaNa}) \setminus \mathbb{N}(\text{EaNa}) : \sim \lambda \lambda A \lambda B \text{Past}(\text{"create } A \ B)$ 
darkness :  $\square(\text{CNs}(\text{n}) \setminus \text{Nt}(\text{s}(\text{n}))) : \sim (\text{"darkness"}, \text{gen } \text{"darkness"})$ 
deep :  $\square \text{CNs}(\text{n}) : \text{deep}$ 
did :  $\mathbb{N}YgVgVYt(\exists ((\exists \text{Na} \ S g) \setminus (\exists \text{Na} \ S \text{b})) \setminus (\text{E} \setminus \text{C} \setminus \text{N} \setminus \text{C} \setminus \text{S} \setminus \text{f})) \setminus ((\exists \text{Na} \ S g) \setminus (\exists \text{Na} \ S \text{b})) : \lambda \lambda A \lambda B((A \ B) \ B)$ 
did+to :  $\square((\exists \text{Na} \ S \text{b}) \setminus (\text{N} \ \text{C} \ \text{S} \ \text{D})) \setminus (\text{N} \ \text{C} \ \text{S} \ \text{D}) \setminus ((\exists \text{Na} \ S \text{f}) \setminus ((\exists \text{Na} \ S \text{g}) \setminus \text{N} \ \text{I} \setminus \text{S} \ \text{f})) : \lambda \text{K} \ \lambda \text{L} \setminus (\text{K} \ \text{L} \ \text{L})$ 
doesnt :  $\mathbb{N}YgVta(\text{Sg} \setminus ((\exists \text{Na} \ S f) \setminus (\exists \text{Na} \ S \text{b}))) \setminus Sg : \lambda A \neg (A \ \lambda B \lambda C(B \ C))$ 
dog :  $\square \text{CNs}(\text{n}) : \text{dog}$ 
donuts :  $\square(\text{Nt}(\text{pr}) \setminus \text{CNpr}(\text{n})) : \sim (\text{gen } \text{"donuts"}, \text{"donuts"})$ 
earth :  $\square \text{CNs}(\text{n}) : \text{earth}$ 
eat :  $\square((\exists \text{EaNa} \ S \text{b}) \ \text{EaNa}) \setminus \mathbb{N}(\text{EaNa}) : \sim \lambda \lambda A \lambda B(\text{"eat } A \ B)$ 
edinburgh :  $\mathbb{N}Ht(s(m)) : e$ 
editor :  $\square(\exists \text{CNs}(\text{g}) \setminus \text{PPof}) : \text{editor}$ 
every :  $\mathbb{N}YgVf((Sf \setminus \text{Nt}(s(g)))Sf) \setminus \text{CNs}(\text{g}) : \lambda \lambda A \lambda B \forall C[(A \ C) \rightarrow (B \ C)]$ 

```

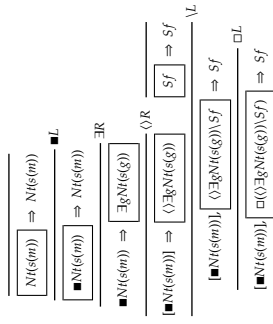
everyone : $\square \forall f((S \uparrow \forall g \text{Nt}(g)) \rightarrow S f) : \sim \lambda \forall B [(\text{"person } B \rightarrow (A \ B))]$
face : $\square \text{CN}(s(n)) : \text{face}$
fell : $\square (\exists n)(\lambda n \setminus S f) : \sim \lambda A (\text{Past } C \text{fall } A)$
filed : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) / \exists n A n : \sim \lambda \lambda A B (\text{Past } ("file A \ B))$
finds : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) / \exists n A n : \sim \lambda \lambda A B (\text{Pres } ("find A \ B))$
fish : $\square \text{CN}(s(n)) : \text{fish}$
for : $\blacksquare (\text{P} / \text{for} / \exists n A n) : \lambda A A$
form : $\square \text{CN}(s(n) \& \text{Nt}(s(n))) : \sim C \text{form, (gen "form)}$
fortunately : $\square \forall f (S \uparrow S f) : \text{fortunately}$
friends : $\square (\text{CN} / \text{P} / \text{of}) : \text{friends}$
from : $\square ((\forall a \forall f ((\lambda n A (S f)) (\lambda n A (S f)) \& \forall n (\text{CN} \setminus \text{CN} n)) / \exists n B n) : \sim \lambda A ("from \text{to } A), (C \text{from} \text{to} A))$
gave : $\square (\square (\exists n A (S f)) / \exists n A n) : \sim \lambda \lambda A B (\text{Past } ((\text{give } \pi_2 A) \pi_1 A \ B))$
gave : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) / \exists n A n \bullet \text{W}[\text{lie, cold, shoulder}] : \sim \lambda \lambda A B (\text{Past } ("shun A \ B))$
gave : $\square (\square (\exists n A (S f)) / \exists n A n) : \sim \lambda \lambda A B \lambda C (\text{Past } ((\text{give } A \ B) \ C))$
girl : $\square \text{CN}(s(f)) : \text{girl}$
gives : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) / \exists n A n \bullet \text{W}[\text{lie, cold, shoulder}] : \sim \lambda \lambda A B (\text{Pres } ("shun A \ B))$
God : $\blacksquare \text{Nt}(s(m)) : \text{God}$
good : $\square \forall n (\text{CN} / \text{CN} n) : \text{good}$
has : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) / \exists n A n : \sim \lambda \lambda A B (\text{Pres } ("have A \ B))$
he : $\blacksquare \square \sim \forall g (\blacksquare \text{Sg} \blacksquare \text{Nt}(s(m)) / (\square \text{Nt}(s(m))) \setminus S g) : \lambda A A$
heaven : $\square \text{CN}(s(n)) : \text{heaven}$
her : $\blacksquare \forall g \forall a ((\lambda n A (S g)) \blacksquare \text{Nt}(s(f))) \blacksquare ((\lambda n A (S g)) \blacksquare \text{Nt}(s(f)))) : \lambda A A$
himself : $\blacksquare \forall f ((\square \text{Nt}(s(m))) \setminus S f) \text{Nt}(s(m)) \blacksquare (\square \text{Nt}(s(m))) \setminus S f) : \sim \lambda \lambda A B ((A \ B) \ B)$
horse : $\square \text{CN}(s(n)) : \text{horse}$
in : $\square (\forall a \forall f ((\lambda n A (S f)) (\lambda n A (S f))) / \exists n A n : \sim \lambda \lambda A B \lambda C ("in A) (B \ C))$
in : $\square (\forall f (S \uparrow S f) / \exists n A n) : \text{in}$
is : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) / \exists n A n \bullet \text{Eg} (\text{CN} g / \text{CN} g) \setminus (\text{CN} g / \text{CN} g) \rightarrow \text{I}) : \sim \lambda \lambda A B (\text{Pres } (A \rightarrow C) \setminus [B = C] \setminus D : (D \ \lambda E [E = B]) \ B))$
it : $\blacksquare \text{W}[\text{it}] : 0$
it : $\blacksquare \forall f \forall a ((\lambda n A (S f)) \blacksquare \text{Nt}(s(n))) \blacksquare ((\lambda n A (S f)) \blacksquare \text{Nt}(s(n)))) : \lambda A A$
it : $\blacksquare \square \sim \forall f (\blacksquare \text{S} \blacksquare \text{Nt}(s(m)) / (\square \text{Nt}(s(m))) \setminus S f) : \lambda A A$
jogs : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) : \sim \lambda A (\text{Pres } ("jog A))$
john : $\blacksquare \text{Nt}(s(m)) : j$
laughs : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) : \sim \lambda A (\text{Pres } ("laugh A))$
left : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) : \sim \lambda A (\text{Pres } ("leave A))$
let : $\square (\text{Simi} / \text{Sb}) : \text{let}$
light : $\square (\text{CN}(s(n) \& \text{Nt}(s(n))) : \sim C \text{light, (gen "light)}$
likes : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) / \exists n A n : \sim \lambda \lambda A B (\text{Pres } ("like A \ B))$
logic : $\square (\text{Nt}(s(n)) \& \text{CN}(s(m)) : \sim (\text{gen "logic, "logic}$
london : $\blacksquare \text{Nt}(s(m)) : l$
loses : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) / \exists n A n : \sim \lambda \lambda A B (\text{Pres } ("lose A \ B))$
love : $\square (\square (\exists n A (S b) / \exists n A n) : \sim \lambda \lambda A B ("love A \ B)$
loved : $\square \forall a \forall b ((\lambda n A (S \rightarrow) \text{Nt}(b)) (\square (\lambda n A (S \rightarrow) \text{Nt}(b)) \blacksquare g (\text{CN} g \setminus \text{CN} g))) : \sim C \text{love, } \lambda \lambda \lambda A B \lambda C [(B \ C) \wedge \exists D ((A \ C) \ D)])$
loves : $\square (\square (\exists g \text{Nt}(s(g))) \setminus S f) / \exists n A n : \sim \lambda \lambda A B (\text{Pres } ("love A \ B))$
man : $\square \text{CN}(s(m)) : \text{man}$
mary : $\blacksquare \text{Nt}(s(f)) : m$
met : $\square (\square (\exists n A (S f) / \exists n A n) : \sim \lambda \lambda A B (\text{Past } ("meet A \ B))$
more : $\blacksquare \forall f \forall g \forall i ((S f) (\text{Sh} \setminus \text{Nt}(p(g))) \setminus S h) / \text{CN} p(g)) \setminus S i / \text{CP} \text{hmm} \blacksquare ((\text{Sh} \setminus \text{Nt}(p(g))) \setminus S h) / \text{CN} p(g))$
mountain : $\square \text{CN}(s(n)) : \text{mountain}$
moved : $\square (\square (\exists n A (S f) : \sim \lambda A (\text{Past } ("move A))$
necessarily : $\blacksquare (\text{SA} / \square \text{SA}) : \text{Nec}$
of : $\square (\forall n (\text{CN} n / \text{CN} n) / \blacksquare \text{Eb} \text{Nt}(b) \& (\text{P} / \text{of} / \exists n A n)) : \sim C \text{of, } \lambda A A$
or : $\blacksquare \forall f ((\blacksquare \text{S} \setminus \square \sim \square \sim S f) / \blacksquare \blacksquare S f) : (\Phi)^{n+0} \text{ or}$

or : $\blacksquare \forall a \forall f ((\exists (\lambda Na) (Sf) \wedge \neg \Pi^{-1} (\lambda (Na) (Sf)) \wedge \neg \Pi^{-1} (\lambda (Na) (Sf))) / \blacksquare (\exists (Na) (Sf))) : (\Phi^{n+} (s \ 0) \text{ or})$
or : $\blacksquare \forall f ((\exists (\lambda (Sg) (\exists \exists g \text{Nh}(s(g)))) (Sf)) \wedge \neg \Pi^{-1} (\exists \exists g \text{Nh}(s(g)))) / \blacksquare (Sf / (\exists \exists g \text{Nh}(s(g)))) (Sf)) : (\Phi^{n+} (s \ 0) \text{ or})$
or : $\blacksquare \forall a \forall f (((\exists (\lambda Na) (Sf) \wedge \neg \Pi^{-1} (\exists (\lambda Na) (Sf) / \exists \exists Nb) / \exists \exists Nb) \wedge \neg \Pi^{-1} (\exists (\lambda Na) (Sf) / \exists \exists Nb) / \exists \exists Nb)) / \blacksquare ((\exists (\lambda Na) (Sf) / \exists \exists Nb) / \exists \exists Nb)) : (\Phi^{n+} (s \ 0) \text{ or})$
painting : $\square \text{CN}(s(n) / \text{ppof}) : \sim \lambda A ((\text{of } A) \text{ . painting})$
paper : $\square \text{CN}(s(n))$
park : $\square \text{CN}(s(n))$
past : $\square \forall a \forall f ((\exists (\lambda Na) (Sf) \wedge \neg \Pi^{-1} (\lambda (Na) (Sf)) / \exists \exists Nb) : \sim \lambda A \lambda B \lambda C ((\sim \text{past } A) (B \ C))$
perseverance : $\square (\text{Nh}(s(n)) \& \text{CN}(s(n))) : (\text{gen } \sim \text{perseverance}, \sim \text{perseverance})$
peter : $\blacksquare \text{Nh}(s(n)) : p$
phonetics : $\square (\text{Nh}(s(n)) \& \text{CN}(s(n))) : (\text{gen } \sim \text{phonetics}, \sim \text{phonetics})$
praises : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) / \exists \exists Na) : \sim \lambda A \lambda B (\text{Pres } ((\sim \text{praise } A) \ B))$
raced : $\square ((\exists \exists Na) (Sf) : \sim \lambda A (\text{Past } (\sim \text{race } A))$
raced : $\square \forall a \forall b ((\exists (\lambda Na) (S - \neg \Pi) b) \wedge ((\exists (\lambda Na) (S - \neg \Pi) b) \wedge \forall g (\text{CN}g \setminus \text{CN}g))) : (\sim \text{race2}, \lambda A \lambda B \lambda C ((B \ C) \wedge \exists D ((A \ C) \ D)))$
rains : $\square ((\exists \exists W[\text{it}] - \text{Sf}) : \sim \text{Pres } \sim \text{it rains})$
reading : $\square ((\exists \exists Na) (S \text{sp}) / \exists \exists Na) : \sim \lambda A \lambda B ((\sim \text{read } A) \ B)$
robin : $\blacksquare \forall g \text{Nh}(s(g)) : r$
said : $\square ((\exists \exists Na) (Sf) / \text{Sim}) : \sim \lambda A \lambda B (\text{Past } ((\sim \text{say } A) \ B))$
saw : $\square ((\exists \exists Na) (Sf) / \exists \exists Na \text{oc} \text{CP} \text{that}) : \sim \lambda A \lambda B (\text{Past } ((A \rightarrow C (\sim \text{see } C) ; D (\sim \text{see } D)) \ B))$
seeks : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) \wedge \forall a \forall f ((\exists (\lambda Na) (Sf) / \exists \exists Nb) \wedge (\exists Na) (Sf)) : \sim \lambda A \lambda B ((\text{Tries } (\sim \text{find } B)) \ B))$
sees : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) / \exists \exists Na) : \sim \lambda A \lambda B (\text{Pres } ((\sim \text{see } A) \ B))$
sent : $\square ((\exists \exists Na) (Sf) / \exists \exists Nb \bullet \text{P} \text{to}) : \sim \lambda A \lambda B (\text{Past } ((\sim \text{send } \pi_2 A) \ \pi_1 A) \ B))$
sent : $\square ((\exists \exists Na) (Sf) / \exists \exists Na) : \sim \lambda A \lambda B \lambda C (\text{Past } ((\sim \text{send } A) \ B) \ C))$
she : $\blacksquare \Pi^{-1} \forall g (\blacksquare Sg \blacksquare \text{Nh}(s(f)) / (\exists \exists \text{Nh}(s(f))) (Sg)) : \lambda A A$
sings : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) : \sim \lambda A (\text{Pres } (\sim \text{sing } A))$
slept : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) : \sim \lambda A (\text{Past } (\sim \text{sleep } A))$
slowly : $\square \forall a \forall f ((\exists (\lambda Na) (Sf) / \exists \exists Na) (Sf)) : \sim \lambda A \lambda B (\text{Pres } (\sim \text{slowly } (\sim \text{A} \cdot B))$
sneezed : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) : \sim \lambda A (\text{Past } (\sim \text{sneeze } A))$
someone : $\square \forall f (Sf \blacksquare \forall g \text{Nh}(g) \setminus Sf) : \sim \lambda A \lambda B (\text{Pres } (((\sim \text{sell } \pi_2 A) \ \pi_1 A) \ B))$
Spirit : $\square \text{CN}(s(m)) : \text{spirit}$
such+that : $\blacksquare \forall n ((\text{CN}n / \text{CN}m) / (Sf \blacksquare \text{Nh}(n))) : \lambda A \lambda B \lambda C ((B \ C) \wedge (A \ C))$
studies : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) / \exists \exists Na) : \sim \lambda A \lambda B (\text{Pres } ((\sim \text{study } A) \ B))$
such+that : $\blacksquare \forall n ((\text{CN}n / \text{CN}m) / (Sf \blacksquare \text{Nh}(n))) : \lambda A \lambda B \lambda C ((B \ C) \wedge (A \ C))$
suzy : $\blacksquare \text{Nh}(s(f)) : s$
talks : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) : \sim \lambda A (\text{Pres } (\sim \text{talk } A))$
tall : $\square \forall g (\text{CN}g / \text{CN}g) : \text{tall}$
teetotal : $\square \forall n (\text{CN}n / \text{CN}n) : \sim \lambda A \lambda B ((A \ B) \wedge (\sim \text{teetotal } B))$
tenmilliondollars : $\square \text{Nh}(s(n)) : \text{tenmilliondollars}$
than : $\blacksquare (\text{CP} \text{that} / \square Sf) : \lambda A A$
that : $\blacksquare \forall n (\neg \Pi^{-1} (\text{CN}n / \text{CN}n)) / \blacksquare ((\exists \text{Nh}(n)) \setminus Sf) : \lambda A \lambda B \lambda C ((B \ C) \wedge (A \ C))$
the : $\blacksquare \forall n (\text{Nh}(n) / \text{CN}n) : t$
the+cold+shoulder : $\blacksquare W[\text{the, cold, shoulder}] : 0$
there : $\blacksquare W[\text{there}] : 0$
thinks : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) / \text{CP} \text{that} \cup \square Sf) : \sim \lambda A \lambda B (\text{Pres } ((\sim \text{think } A) \ B))$
to : $\blacksquare (P \ \text{P} \ \text{to} / \exists \exists Na) \wedge \forall n ((\exists (\lambda Na) (Sf) / \exists \exists Nb) \wedge (\exists (\lambda Na) (Sf) / \exists \exists Nb)) : \lambda A A$
today : $\square \forall a \forall f ((\exists (\lambda Na) (Sf) \wedge \neg \Pi^{-1} (\lambda (Na) (Sf)) : \sim \lambda A \lambda B ((\sim \text{today } (A \ B))$
tries : $\square ((\exists \exists g \text{Nh}(s(g))) (Sf) / \exists \exists Na) : \sim \lambda A \lambda B ((\sim \text{Tries } (\sim \text{A} \ B)) \ B))$
unicorn : $\square \text{CN}(s(n)) : \text{unicorn}$
up : $\blacksquare W[\text{up}] : 0$
upon : $\square ((\forall a \forall f ((\exists (\lambda Na) (Sf) \wedge \neg \Pi^{-1} (\lambda (Na) (Sf)) \& \forall g (\text{CN}g \setminus \text{CN}g)) / \exists \exists Na) : \sim \lambda A ((\sim \text{upon} \text{ad} A), (\sim \text{upon} \text{ad}n \ A)))$
void : $\square \forall g (\text{CN}g / \text{CN}g) : \text{void}$
walk : $\square ((\exists \exists Na) (Sf) / \exists \exists Na) : \sim \lambda A (\text{Pres } (\sim \text{walk } A))$

walk : $\square(\exists \lambda \exists a \lambda a \backslash S b) : \sim \lambda A ({}^c \text{walk } A)$
walks : $\square(\exists g \exists N f(s(g)) \backslash S f) : \sim \lambda A ({}^c \text{Pres } ({}^c \text{walk } A))$
was : $\blacksquare(\exists g \exists N f(s(g)) \backslash S f) / (\exists a \exists N \text{mem}(\exists g(\text{CN}g / \text{CN}g) \sqcup (\text{CN}g / \text{CN}g) \sqcup (\text{CN}g / \text{CN}g)) \rightarrow I)) : \lambda A \lambda B (\text{Past } (A \rightarrow C [B = C] ; D ; D . (\lambda E [E = B] ; B)))$
was : $\square(\exists W [{}^c \text{here}] \rightarrow S f) / \exists a \lambda a \lambda a : \sim \lambda A (\text{Past } ({}^c \text{be } A))$
waters : $\square \text{CN} p(n) : \text{waters}$
which : $\blacksquare \forall n \forall m ((N f(n) \backslash N f(m)) \sqcup (\neg \text{I} \neg \text{I}^{-1} (\text{CN} m / \text{CN} m)) / \blacksquare((\exists N f(n) \sqcup \blacksquare N f(m)) \backslash S f)) : \lambda A \lambda B \lambda C \lambda D [(C D) \wedge (B (A D))]$
who : $\blacksquare \forall n \forall m (\text{I} \neg \text{I}^{-1} \neg (N f(n) \backslash (S f \backslash N f(m) \backslash S h)) / \blacksquare((\exists N f(n) \sqcup \blacksquare N f(m)) \backslash S f)) : \lambda A \lambda B (\text{Fut } (A B))$
will : $\blacksquare \forall a (\exists N a \backslash S f) / (\exists N a \backslash S b) : \lambda A \lambda B (\text{Fut } (A B))$
without : $\square(\forall g(\text{CN}g / \text{CN}g) / \exists a \lambda a : \sim \lambda A \lambda B \lambda C [(B C) \wedge \neg ({}^c \text{with } A C)]$
without : $\blacksquare \forall a f (\text{I} \neg \text{I}^{-1} ((\exists N a \backslash S f) \backslash ((\exists N a \backslash S f)) \backslash ((\exists N a \backslash S p s p)) : \lambda A \lambda B \lambda C [(B C) \wedge \neg (A C)]$
woman : $\square \text{CN} s(f) : \text{woman}$
yesterday : $\square \forall a \forall f ((\exists N a \backslash S f) \backslash ((\exists N a \backslash S f)) : \sim \lambda A \lambda B ({}^c \text{yesterday } (A B))$

(dwp(7-7)) {john}+walks : Sf

$\blacksquare N f(s(m)) : \text{I} ; \square(\exists g \exists N f(s(g)) \backslash S f) : \sim \lambda A (\text{Pres } ({}^c \text{walk } A)) \Rightarrow S f$



(Pres (walk))

(dwp(7-16)) {every+man}+talks : Sf

$\blacksquare \forall g (\forall f (S f \backslash N f(s(g))) \backslash S f) / \text{CN}(s(g)) : \lambda A \lambda B \lambda C [(A C) \rightarrow (B C)] ; \square \text{CN}(s(m)) : \text{man} ; \square(\exists g \exists N f(s(g)) \backslash S f) : \sim \lambda D (\text{Pres } ({}^c \text{talk } D)) \Rightarrow S f$

(Pres (walk (r_fish)))

(dwp(Vf(Sf(Nt(e(g))))Sf))CN(e(g)) : $\lambda \lambda \text{BVC}(A \rightarrow B \text{ C}), \text{CN}(s(m) : \text{mm}), [\Box(\exists \text{EgNt}(e(g))Sf) : \neg \text{AD}(\text{Pres}(\text{walk} D)), \blacksquare \text{Vf}(\Box \text{Sf}) \Pi^{-1} \Pi^{-1} \text{Sf}] \blacksquare \text{Sf} : \text{e}^{\text{pr}^+} \text{ o on}, \Box(\exists \text{EgNt}(e(g))Sf) : \neg \text{LE}(\text{Pres}(\text{walk} E))] \Rightarrow \text{Sf}$

$\blacksquare \text{EgVf}(\text{Sf}^{\uparrow} \text{Nt}(e(g)))Sf) \text{CN}(e(g)) : \lambda \lambda \text{BVC}(A \rightarrow B \text{ C}), \text{CN}(s(m) : \text{mm}), [\Box(\exists \text{EgNt}(e(g))Sf) : \neg \text{AD}(\text{Pres}(\text{walk} D)), \blacksquare \text{Vf}(\Box \text{Sf}) \Pi^{-1} \Pi^{-1} \text{Sf}] \blacksquare \text{Sf} : \text{e}^{\text{pr}^+} \text{ o on}, \Box(\exists \text{EgNt}(e(g))Sf) : \neg \text{LE}(\text{Pres}(\text{walk} E))] \Rightarrow \text{Sf}$

$\blacksquare \text{EgVf}(\text{Sf}^{\uparrow} \text{Nt}(e(g)))Sf) \text{CN}(e(g)) : \lambda \lambda \text{BVC}(A \rightarrow B \text{ C}), \text{CN}(s(m) : \text{mm}), [\Box(\exists \text{EgNt}(e(g))Sf) : \neg \text{AD}(\text{Pres}(\text{walk} D)), \blacksquare \text{Vf}(\Box \text{Sf}) \Pi^{-1} \Pi^{-1} (\Box \text{NaSf}) \blacksquare (\Box \text{NaSf}) : \text{e}^{\text{pr}^+} (s \text{ o on}), \Box(\exists \text{EgNt}(e(g))Sf) : \neg \text{LE}(\text{Pres}(\text{walk} E))] \Rightarrow \text{Sf}$

$\frac{\text{Nt}(e(m)) \Rightarrow \text{Nt}(e(m))}{\text{Nt}(e(m)) \Rightarrow \text{EgNt}(e(g))} \text{ER}$ $\frac{\text{Nt}(e(m)) \Rightarrow \text{EgNt}(e(g))}{[\text{Nt}(e(m)) \Rightarrow \Box \text{EgNt}(e(g))] \text{Sf} \Rightarrow \text{Sf}} \text{OR}$ $\frac{[\text{Nt}(e(m)) \blacksquare \Box \text{EgNt}(e(g))] \text{Sf} \Rightarrow \text{Sf}}{[\text{Nt}(e(m)) \blacksquare \Box(\exists \text{EgNt}(e(g))Sf)] \Rightarrow \text{Sf}} \text{OL}$ $\frac{[\text{Nt}(e(m)) \blacksquare \Box(\exists \text{EgNt}(e(g))Sf)] \Rightarrow \text{Sf}}{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \text{Nt}(e(m))} \text{OR}$ $\frac{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \text{Nt}(e(m)) \text{Sf}}{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \blacksquare (\Box \text{Nt}(e(m))Sf)} \text{R}$ $\frac{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \blacksquare (\Box \text{Nt}(e(m))Sf)}{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \blacksquare (\Box \text{Nt}(e(m))Sf)} \text{R}$ $\frac{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \blacksquare (\Box \text{Nt}(e(m))Sf)}{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \blacksquare (\Box \text{Nt}(e(m))Sf)} \text{R}$ $\frac{[\Box(\exists \text{EgNt}(e(g))Sf) \blacksquare \blacksquare (\Box \text{Nt}(e(m))Sf)] \Rightarrow \text{Sf}}{[\Box(\exists \text{EgNt}(e(g))Sf) \blacksquare \blacksquare (\Box \text{Nt}(e(m))Sf)] \Rightarrow \text{Sf}} \text{IL}$	$\frac{\text{Nt}(e(m)) \Rightarrow \text{Nt}(e(m))}{\text{Nt}(e(m)) \Rightarrow \text{EgNt}(e(g))} \text{ER}$ $\frac{\text{Nt}(e(m)) \Rightarrow \text{EgNt}(e(g))}{[\text{Nt}(e(m)) \Rightarrow \Box \text{EgNt}(e(g))] \text{Sf} \Rightarrow \text{Sf}} \text{OR}$ $\frac{[\text{Nt}(e(m)) \blacksquare \Box \text{EgNt}(e(g))] \text{Sf} \Rightarrow \text{Sf}}{[\text{Nt}(e(m)) \blacksquare \Box(\exists \text{EgNt}(e(g))Sf)] \Rightarrow \text{Sf}} \text{OL}$ $\frac{[\text{Nt}(e(m)) \blacksquare \Box(\exists \text{EgNt}(e(g))Sf)] \Rightarrow \text{Sf}}{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \text{Nt}(e(m))} \text{OR}$ $\frac{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \text{Nt}(e(m)) \text{Sf}}{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \blacksquare (\Box \text{Nt}(e(m))Sf)} \text{R}$ $\frac{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \blacksquare (\Box \text{Nt}(e(m))Sf)}{\Box(\exists \text{EgNt}(e(g))Sf) \Rightarrow \blacksquare (\Box \text{Nt}(e(m))Sf)} \text{R}$ $\frac{[\Box(\exists \text{EgNt}(e(g))Sf) \blacksquare \blacksquare (\Box \text{Nt}(e(m))Sf)] \Rightarrow \text{Sf}}{[\Box(\exists \text{EgNt}(e(g))Sf) \blacksquare \blacksquare (\Box \text{Nt}(e(m))Sf)] \Rightarrow \text{Sf}} \text{IL}$
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[[\blacksquare vsV7((sf[↑]Ni(s(g)))↓Sf)/CNs(g) : λ AABVC(A C) → (B C), \square CN_{s(m)} : man , \square (\exists gN(s(g))\Sf) : λ LD(Pres ('talk D)), \blacksquare vsv7((\square (\square Na Sf))\blacksquare(O_{Na}Sf))\blacksquare(O_{Na}Sf) : (ϕ ⁺ (s 0) or), \blacksquare vsv7((sf[↑]Ni(s(g)))↓Sf)/CNs(g) : λ E1FVC(E C) → (F G), \square CN_{s(m)} : man , \square (\exists gN(s(g))\Sf) : \Rightarrow Sf

[[\blacksquare vsV7((sf[↑]Ni(s(g)))↓Sf)/CNs(g) : λ AABVC(A G) → (B C), \square CN_{s(m)} : man , \square (\exists gN(s(g))\Sf) : λ H(Pres ('talk H))]] \Rightarrow Sf

[[\blacksquare vsV7((sf[↑]Ni(s(g)))↓Sf)/CNs(g) : λ AABVC(A G) → (B C), \square CN_{s(m)} : man , \square (\exists gN(s(g))\Sf) : λ LD(Pres ('talk D)), \blacksquare vsv7((\square (\square Na Sf))\blacksquare(O_{Na}Sf))\blacksquare(O_{Na}Sf) : (ϕ ⁺ (s 0) or), \blacksquare vsv7((sf[↑]Ni(s(g)))↓Sf)/CNs(g) : λ E1FVC(E G) → (F G), \square CN_{s(m)} : man , \square (\exists gN(s(g))\Sf) : \Rightarrow Sf

[[\blacksquare vsV7((sf[↑]Ni(s(g)))↓Sf)/CNs(g) : λ AABVC(A G) → (B C), \square CN_{s(m)} : man , \square (\exists gN(s(g))\Sf) : λ LD(Pres ('talk D)), \blacksquare vsv7((\square (\square Na Sf))\blacksquare(O_{Na}Sf))\blacksquare(O_{Na}Sf) : (ϕ ⁺ (s 0) or), \blacksquare vsv7((sf[↑]Ni(s(g)))↓Sf)/CNs(g) : λ E1FVC(E G) → (F G), \square CN_{s(m)} : man , \square (\exists gN(s(g))\Sf) : λ H(Pres ('talk H))]] \Rightarrow Sf

(dwp(7-39)) [[[a+woman]+walks+and+[she]+talks]] : Sf

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